

Interference Avoidance for Wireless Networks

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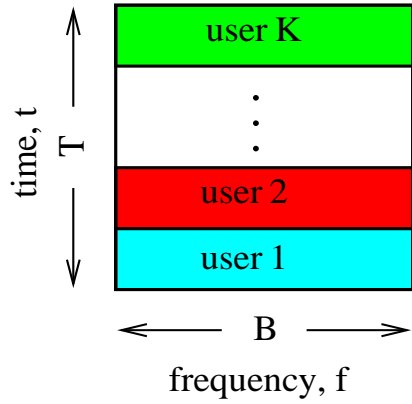
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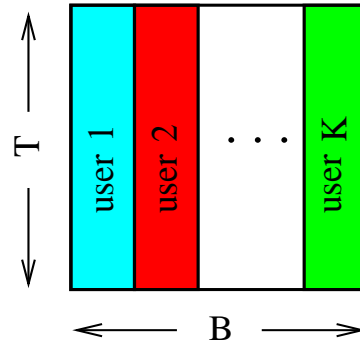
Time Division
Multiple Access

TDMA



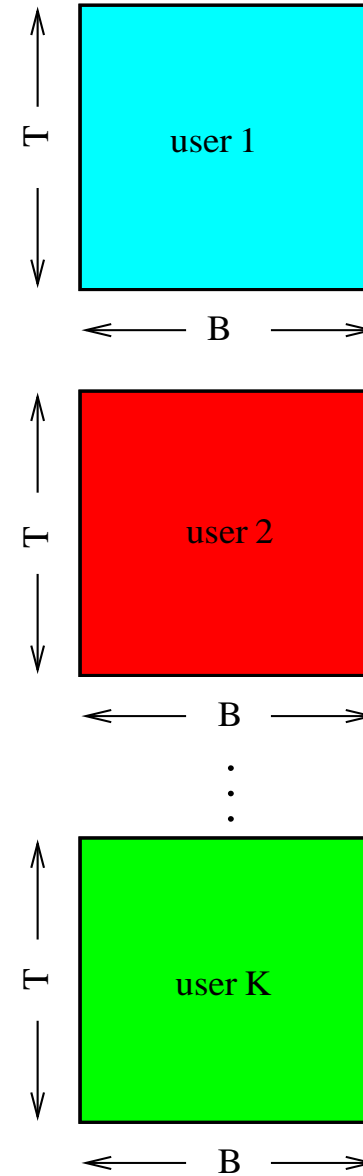
Frequency Division
Multiple Access

FDMA



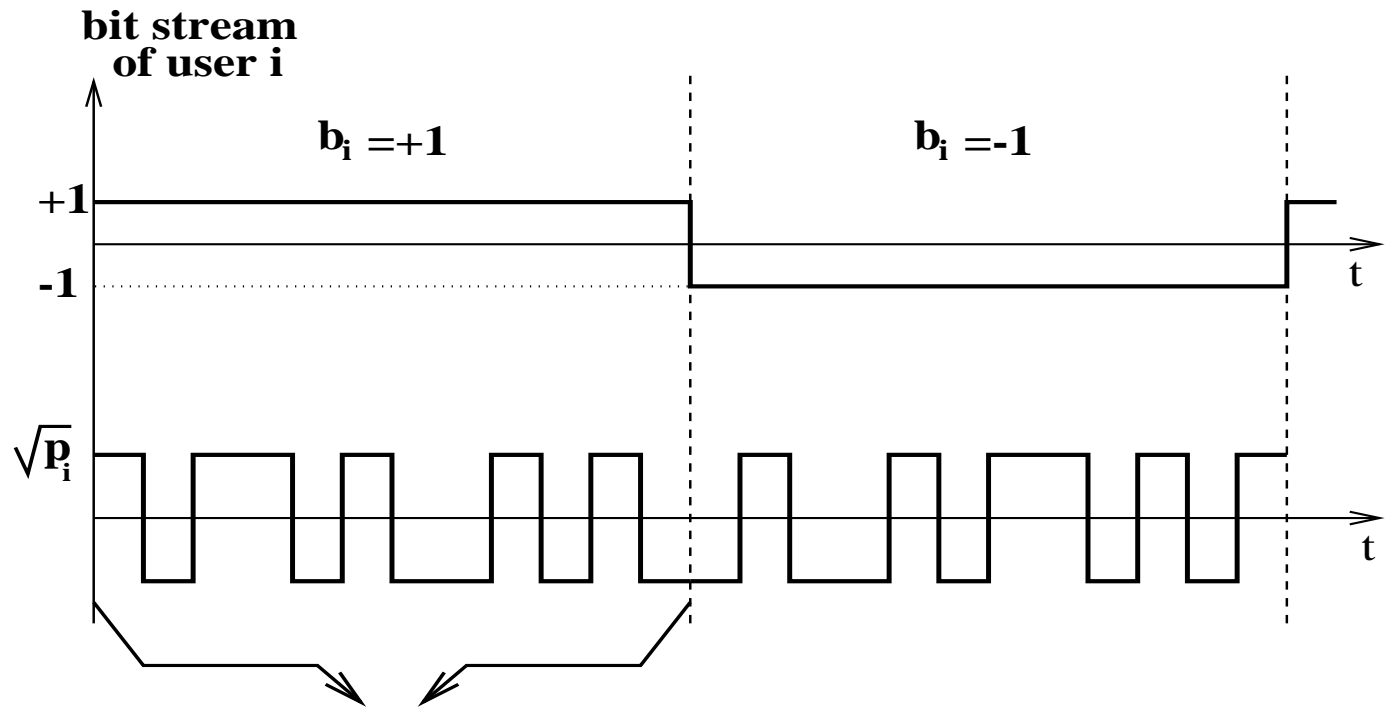
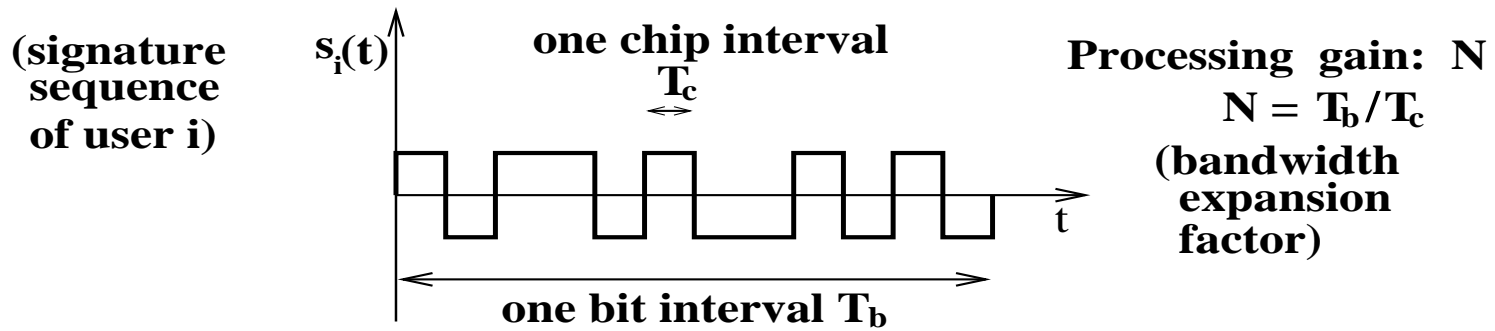
Code Division
Multiple Access

CDMA



**MULTIPLE ACCESS
METHODS**

CDMA Principles



in one bit interval = $\sqrt{p_i} b_i s_i(t)$

From Continuous Signals to Signal Vectors

- Received signal in a given bit interval at the base station of user i is

$$r(t) = \sum_{j=1}^K \sqrt{p_j h_{ij}} b_j s_j(t) + n(t)$$

- Chip matched filtered and sampled signal

$$\mathbf{r} = \sum_{j=1}^K \sqrt{p_j h_{ij}} b_j \mathbf{s}_j + \mathbf{n}$$

- Detection problem:

given the received signal \mathbf{r} , find out what was transmitted by each user, $\{b_1, \dots, b_K\}$

Optimum Multiuser Detection

- Define the bit vector $\mathbf{b} = [b_1, \dots, b_K]^\top$
- For equally likely messages, choose bits of the users according to

$$\max_{\mathbf{b}} f(\mathbf{r} | \mathbf{b})$$

- Jointly optimum multiuser detection (Verdú, 1984)

$$\min_{\mathbf{b} \in \{\pm 1\}^K} \mathbf{b}^\top \mathbf{H} \mathbf{b} - 2\mathbf{b}^\top \mathbf{A} \mathbf{y}$$

where $\mathbf{H} = \mathbf{A} \mathbf{R} \mathbf{A}$,

\mathbf{R} is the cross correlation matrix, $R_{ij} = \mathbf{s}_i^\top \mathbf{s}_j$, and

\mathbf{A} is the (diagonal) amplitude matrix, $A_{ii} = \sqrt{p_i h_{ii}}$

- Nice convex cost function, but discrete feasible set.
- Need to check all possible 2^K bit combinations.
- The optimum multiuser detection is **NP-hard** (has **exponential** computational complexity)

Linear Multiuser Detection

- Given the received signal,

$$\mathbf{r} = \sum_{j=1}^K \sqrt{p_j h_{ij}} b_j \mathbf{s}_j + \mathbf{n}$$

choose \mathbf{c}_i and make bit detection according to

$$\hat{b}_i = \text{sign}(\mathbf{c}_i^\top \mathbf{r})$$

- Single-user matched filter:** $\mathbf{c}_i = \mathbf{s}_i$
- Decorrelator:** Choose \mathbf{c}_i to $\max(\mathbf{c}_i^\top \mathbf{s}_i)^2$ while $\mathbf{c}_i^\top \mathbf{s}_j = 0$ for $j \neq i$
- MMSE receiver:** Choose \mathbf{c}_i to $\min E [(\mathbf{c}_i^\top \mathbf{r} - b_i)^2]$

$$\begin{aligned} \mathbf{c}_i &= \alpha_i \left(\sum_{j=1}^K p_j h_{ij} \mathbf{s}_j \mathbf{s}_j^\top + \sigma^2 \mathbf{I}_N \right)^{-1} \mathbf{s}_i \\ &= \alpha_i \left(\mathbf{S} \mathbf{A}^2 \mathbf{S}^\top + \sigma^2 \mathbf{I}_N \right)^{-1} \mathbf{s}_i \end{aligned}$$

Near-Far Problem

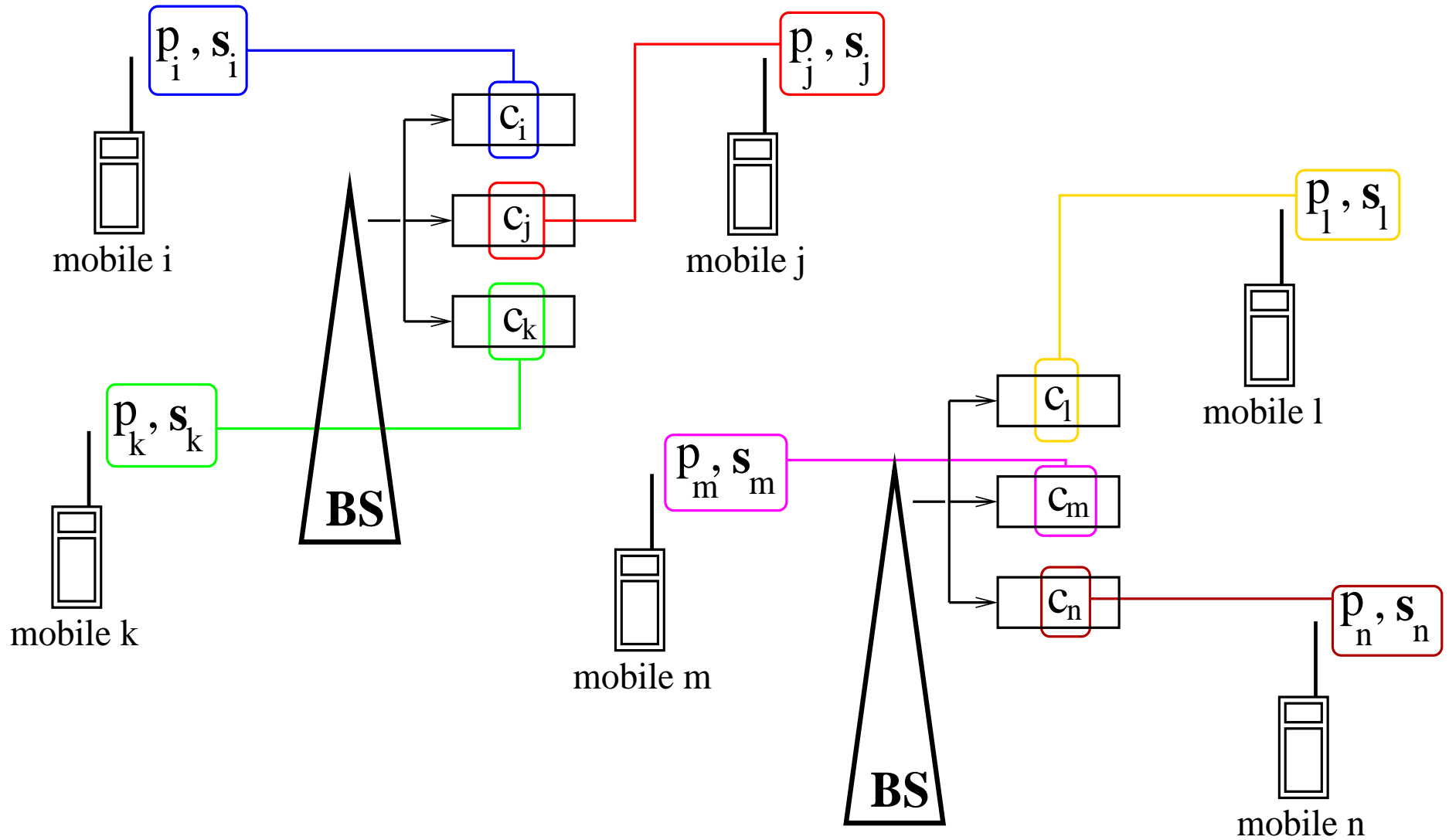
- For the i th user

$$y_i = \mathbf{c}_i^\top \mathbf{r} = \underbrace{\sqrt{p_i h_{ii}} b_i(\mathbf{c}_i^\top \mathbf{s}_i)}_{\text{desired signal}} + \underbrace{\sum_{j \neq i} \sqrt{p_j h_{ij}} b_j(\mathbf{c}_i^\top \mathbf{s}_j)}_{\text{multiaccess interference}} + \underbrace{(\mathbf{c}_i^\top \mathbf{n})}_{\text{Gaussian}}$$

- Users with high received power can degrade the quality of communication of weak users.
- Possible solutions:
 - **multiuser detection**: suppress the multiaccess interference and/or Gaussian noise
 - **power control**: balance the powers
 - **signature sequence design**: avoid the multiaccess interference
 - **receiver beamforming**: multiple antennas for diversity and interference suppression

Controllable Parameters

- Each user has a power (p_i), a signature sequence (s_i) and a receiver filter (c_i)



SIR Constraints

- Filter output of user i has signal and interference components

$$y_i = \underbrace{\sqrt{p_i h_{ii}} (\mathbf{c}_i^\top \mathbf{s}_i) b_i}_{\text{desired signal}} + \underbrace{\sum_{j \neq i} \sqrt{p_j h_{ij}} (\mathbf{c}_i^\top \mathbf{s}_j) b_j + (\mathbf{c}_i^\top \mathbf{n}_i)}_{\text{interference}}$$

- Link quality is determined by the SIR of the user

$$\text{SIR}_i = \frac{\text{Signal Power}}{\text{Interference Power}} = \frac{p_i h_{ii} (\mathbf{c}_i^\top \mathbf{s}_i)^2}{\sum_{j \neq i} p_j h_{ij} (\mathbf{c}_i^\top \mathbf{s}_j)^2 + \sigma^2 (\mathbf{c}_i^\top \mathbf{c}_i)}$$

- β_i : SIR target of user i .
- SIR target has a 1-to-1 relationship with the Bit Error Rate (BER)

$$\text{BER}_i \leq \varepsilon_i \quad \iff \quad \text{SIR}_i \geq \beta_i$$

Reduction to Conventional Power Control

- In the general, SIR of user i :

$$\text{SIR}_i = \frac{p_i h_{ii} (\mathbf{c}_i^\top \mathbf{s}_i)^2}{\sum_{j \neq i} p_j h_{ij} (\mathbf{c}_i^\top \mathbf{s}_j)^2 + \sigma^2 (\mathbf{c}_i^\top \mathbf{c}_i)}$$

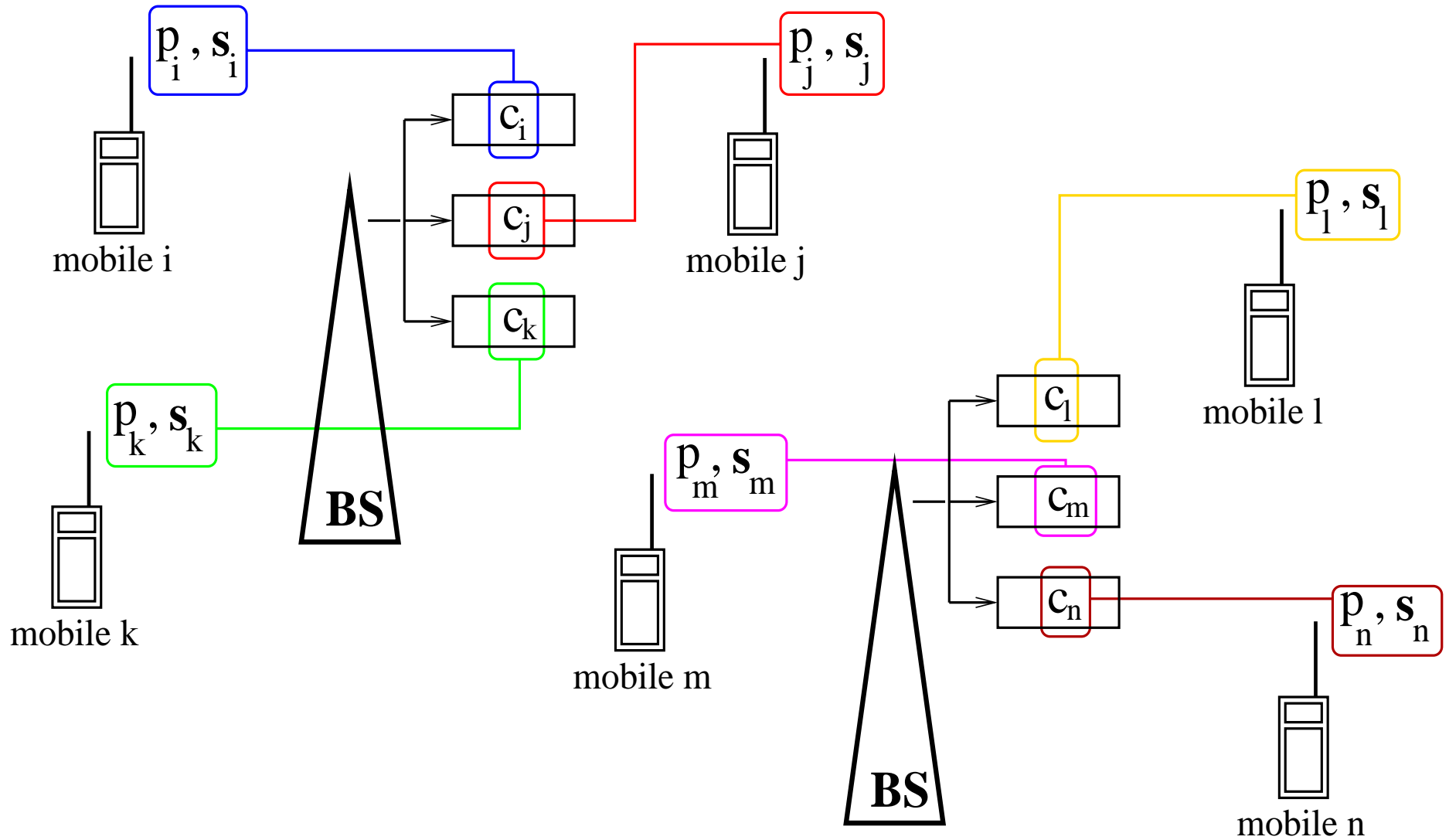
- Fix signature sequences of the users
- Then, let $\mathbf{c}_i = \mathbf{s}_i$, i.e., matched filters
- Note $\mathbf{s}_i^\top \mathbf{s}_i = 1$ and define $\rho_{ij} = \mathbf{s}_i^\top \mathbf{s}_j$

$$\text{SIR}_i = \frac{p_i h_{ii}}{\sum_{j \neq i} p_j h_{ij} \rho_{ij}^2 + \sigma^2}$$

- SIR_i is a function of the power vector only

Controllable Parameters

- Each user has a power (p_i) — signature sequences and receiver filters are fixed



The Conventional Power Control Problem

- Find the *componentwise smallest* power vector such that

$$\text{SIR}_i = \frac{p_i h_{ii}}{\sum_{j \neq i} p_j h_{ij} \rho_{ij}^2 + \sigma^2} \geq \beta_i \quad i = 1, \dots, K$$

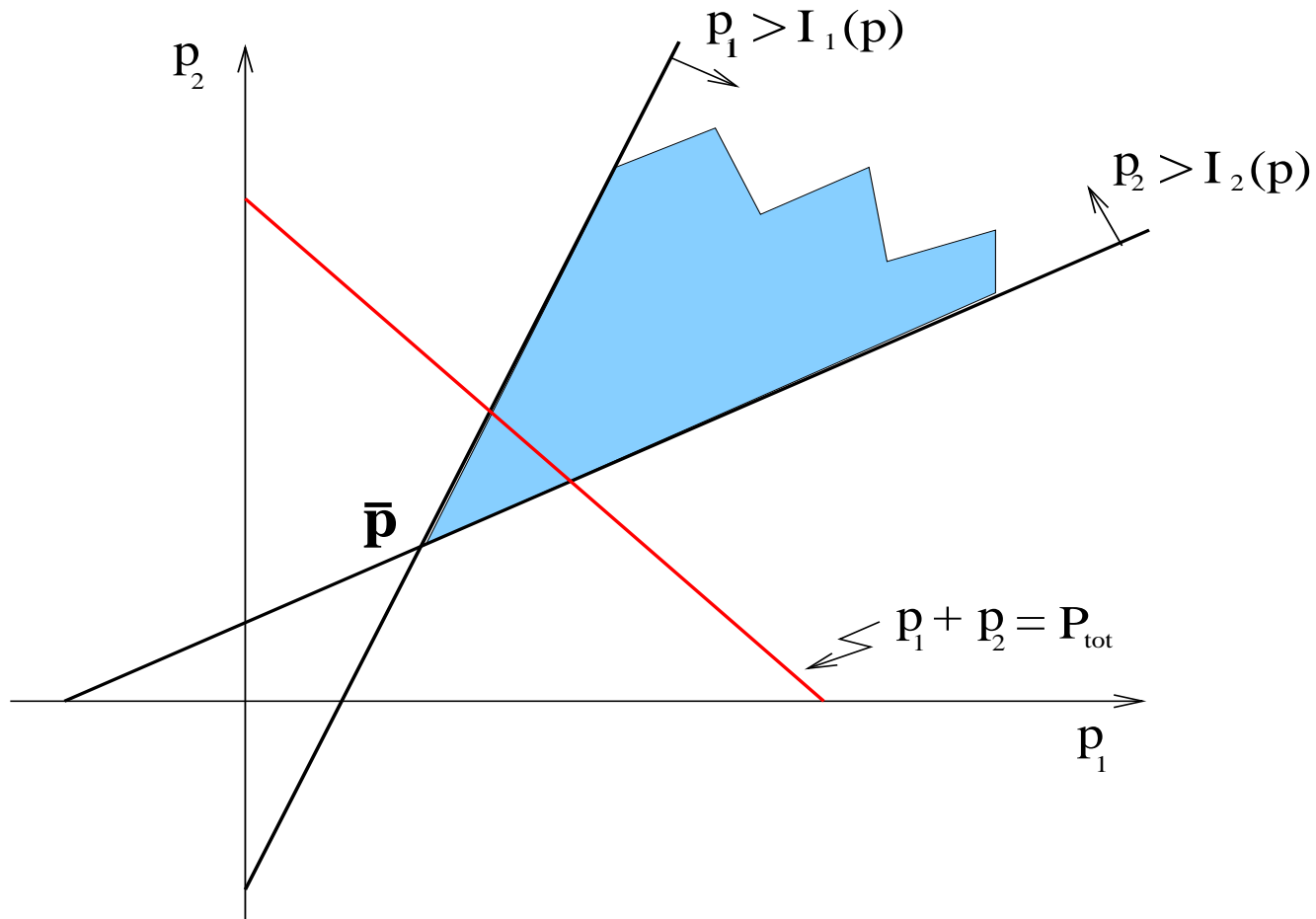
- SIR constraint for the i th user

$$p_i \geq \sum_{j \neq i} \frac{\beta_i h_{ij} \rho_{ij}^2}{h_{ii}} p_j + \frac{\beta_i \sigma^2}{h_{ii}} \quad \iff \quad p_i \geq I_i(\mathbf{p})$$

- In vector notation

$$\mathbf{p} \geq \mathbf{A}\mathbf{p} + \boldsymbol{\eta} \quad \iff \quad \mathbf{p} \geq \mathbf{I}(\mathbf{p})$$

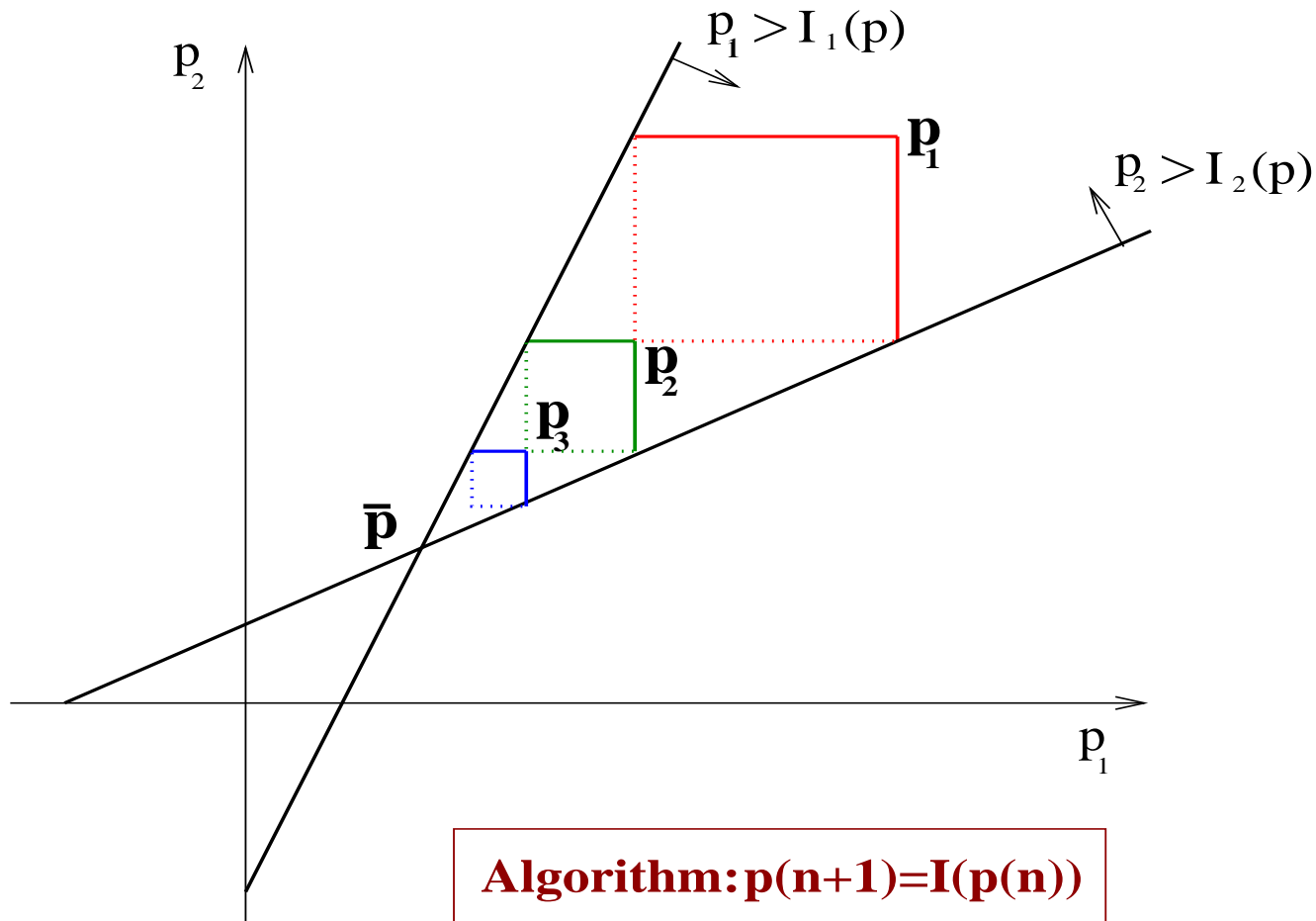
Power Control – Feasible Region



- In general **optimum power vector** is found at $\bar{\mathbf{p}} = \mathbf{I}(\bar{\mathbf{p}})$

$$\bar{\mathbf{p}} = \mathbf{A}\bar{\mathbf{p}} + \boldsymbol{\eta} \quad \Longleftrightarrow \quad \bar{\mathbf{p}} = (\mathbf{I} - \mathbf{A})^{-1} \boldsymbol{\eta}$$

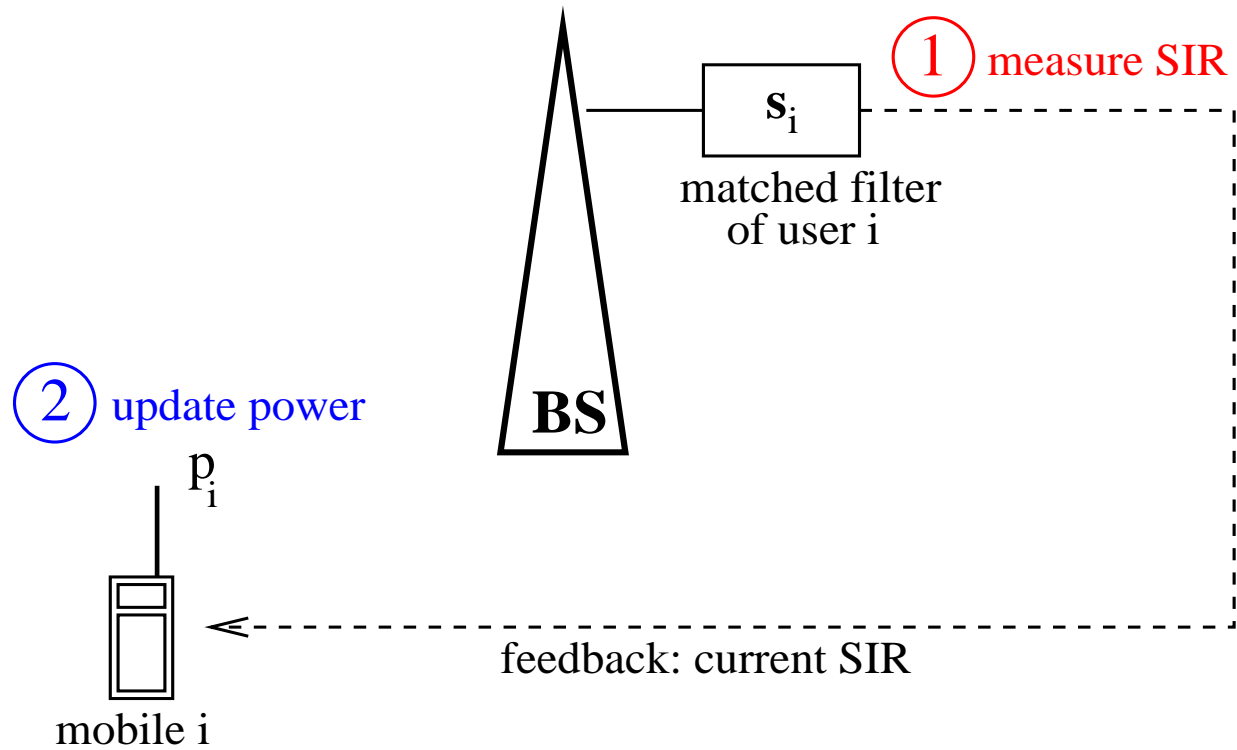
Power Control – A Distributed Algorithm



- **Theorem 1 (Yates)** *If $\mathbf{I}(\mathbf{p})$ is standard (positive, monotone and scalable), then $\mathbf{p}(n+1) = \mathbf{I}(\mathbf{p}(n))$ converges to the componentwise smallest feasible \mathbf{p} .*

Power Control – Implementation

$$\begin{aligned} p_i(n+1) &= I_i(\mathbf{p}(n)) \\ &= \frac{\beta_i}{\text{SIR}_i(n)} p_i(n) \end{aligned}$$



Joint Power Control and Interference Suppression

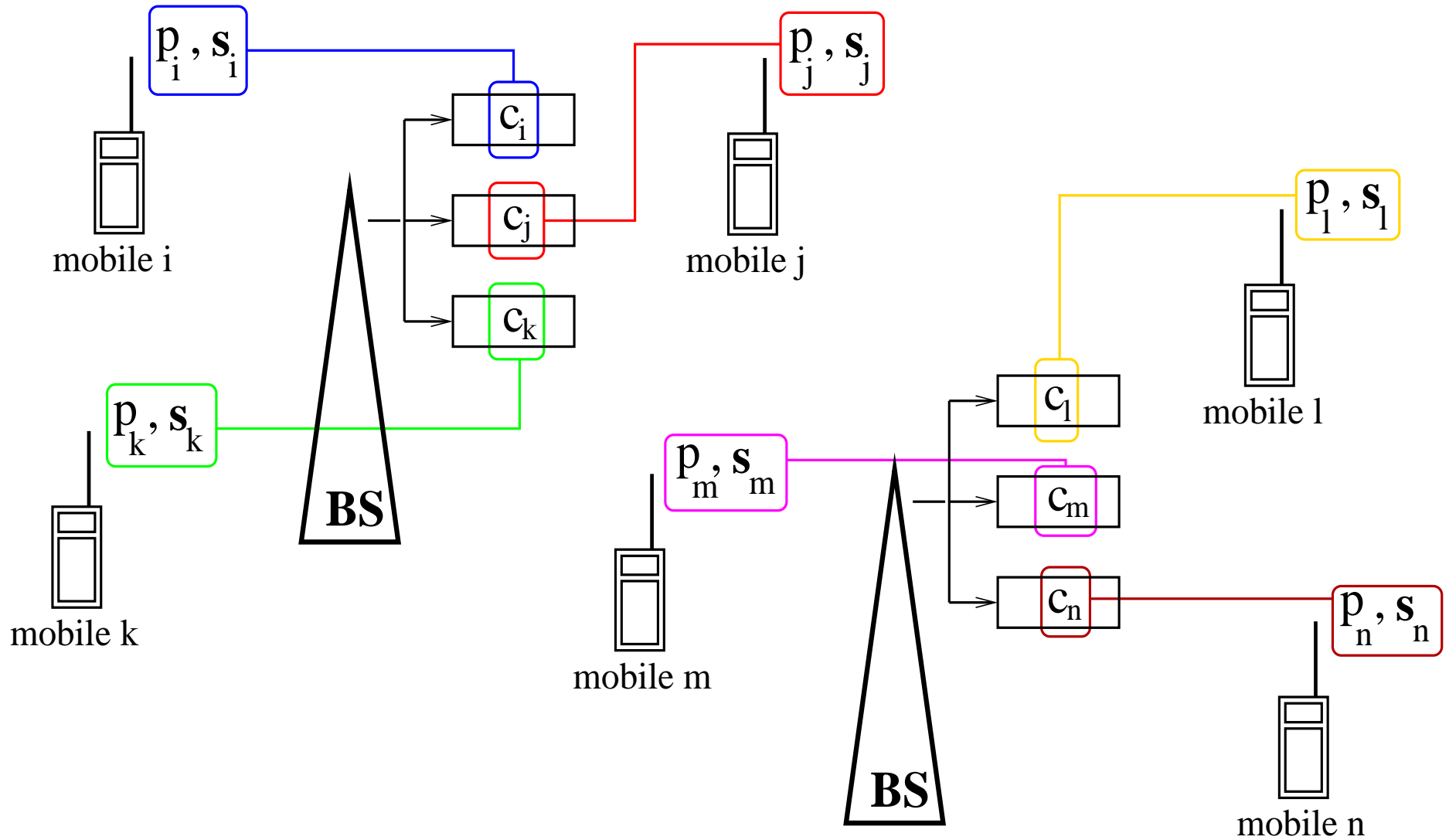
- Fix signature sequences of the users
- But, do not fix \mathbf{c}_i to be the matched filter; treat \mathbf{c}_i as a **free-variable**
- SIR of user i is a function of powers of all users and receiver filter of user i
- Find jointly optimum power vector and receiver filters such that

$$\text{SIR}_i = \frac{p_i h_{ii} (\mathbf{c}_i^\top \mathbf{s}_i)^2}{\sum_{j \neq i} p_j h_{ij} (\mathbf{c}_i^\top \mathbf{s}_j)^2 + \sigma^2 (\mathbf{c}_i^\top \mathbf{c}_i)} \geq \beta_i \quad i = 1, \dots, K$$

with *componentwise smallest* power vector (Ulukus-Yates, 1998)

Controllable Parameters

- Each user has a power (p_i) and a receiver filter (c_i) — signature sequences are fixed



Joint Power Control and Interference Suppression Algorithm

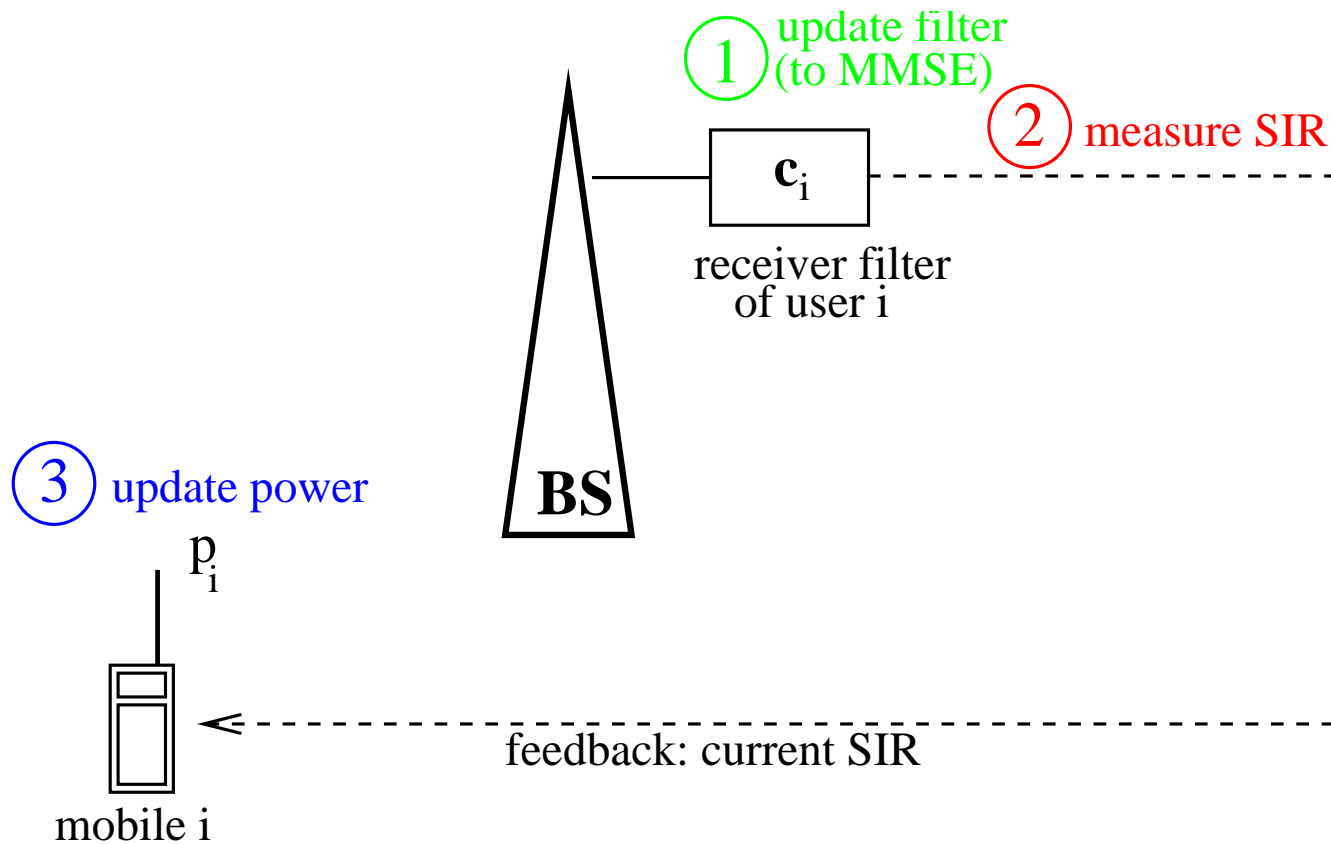
- Define interference function

$$\begin{aligned} I_i(\mathbf{p}) &= \min_{\mathbf{c}_i} \frac{\beta_i}{h_{ii}} \frac{\sum_{j \neq i} p_j h_{ij} (\mathbf{c}_i^\top \mathbf{s}_j)^2 + \sigma^2 (\mathbf{c}_i^\top \mathbf{c}_i)}{(\mathbf{c}_i^\top \mathbf{s}_i)^2} \\ &= \frac{\beta_i}{h_{ii}} \min_{\mathbf{c}_i} \frac{\mathbf{c}_i^\top \left(\sum_{j \neq i} p_j h_{ij} \mathbf{s}_j \mathbf{s}_j^\top + \sigma^2 \mathbf{I}_N \right) \mathbf{c}_i}{(\mathbf{c}_i^\top \mathbf{s}_i)^2} \end{aligned}$$

- **$\mathbf{I}(\mathbf{p})$ is a standard interference function.**
- Power control algorithm: $\mathbf{p}(n+1) = \mathbf{I}(\mathbf{p}(n))$
- The algorithm converges to the componentwise smallest power vector (and corresponding receiver filters) where SIR requirements of the users are satisfied.
- Inherently a two-step algorithm. At each step n ,
 - For fixed transmit powers, find receiver that minimizes the interference: **MMSE filter**
 - Update powers to **just beat** the interference at the output of the updated filters

Joint Power Control and Interference Suppression – Implementation

- Update the receiver filter to be the MMSE filter
- Measure the SIR at the output of the updated filter
- Update power in the usual way: $p_i(n+1) = \frac{\beta_i}{\text{SIR}_i(n)} p_i(n)$.



Joint Power Control, Interference Suppression/Avoidance

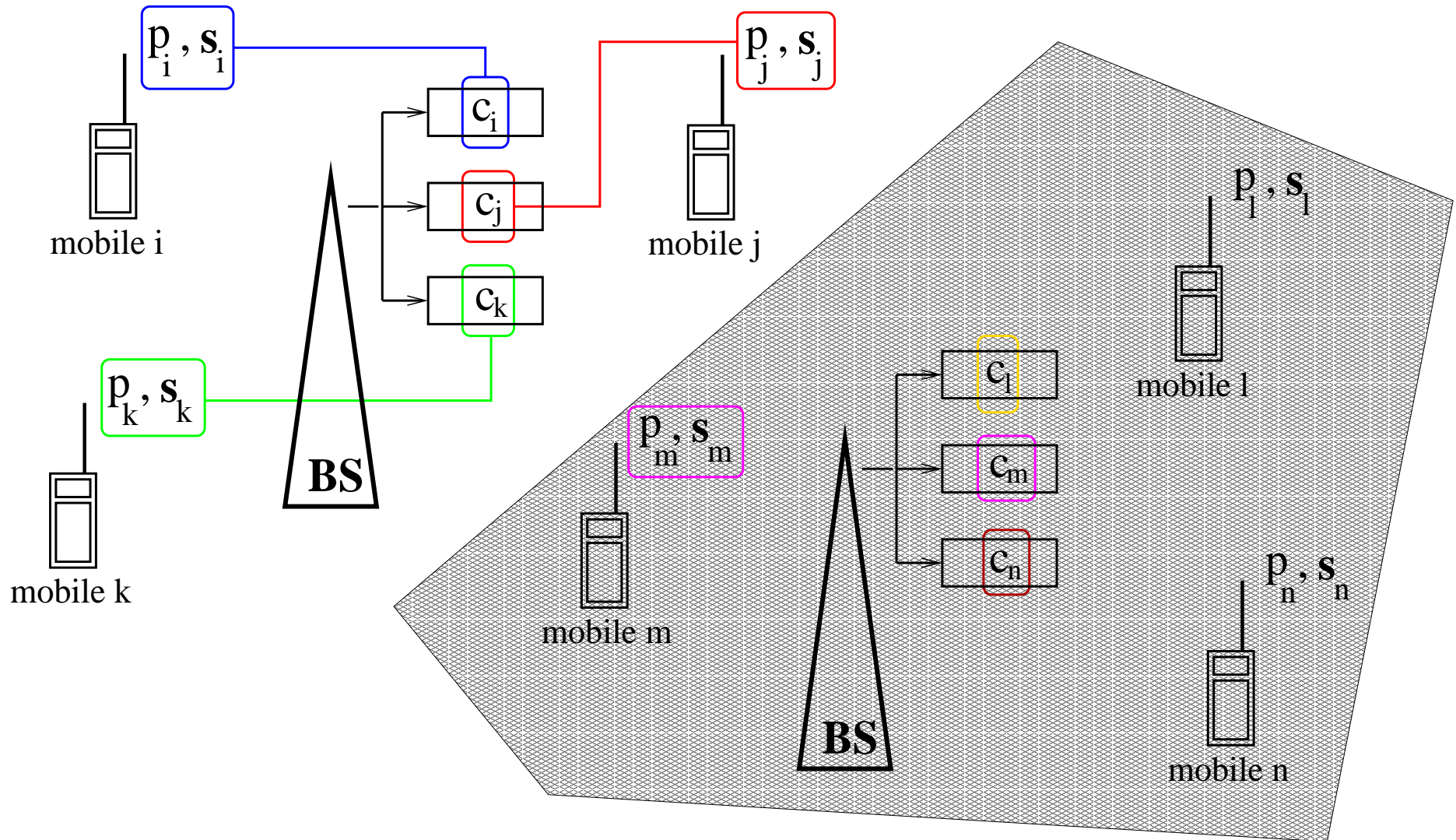
- Do not randomly choose or arbitrarily fix the signature sequences.
- Treat both signature sequences, \mathbf{s}_i , and receiver filters, \mathbf{c}_i , as **free-variables**.
- SIR_i is a function of all powers, all signatures and the receiver filter of user i

$$\text{SIR}_i = \frac{p_i h_{ii} (\mathbf{c}_i^\top \mathbf{s}_i)^2}{\sum_{j \neq i} p_j h_{ij} (\mathbf{c}_i^\top \mathbf{s}_j)^2 + \sigma^2 (\mathbf{c}_i^\top \mathbf{c}_i)}$$

- Choose jointly optimum powers, signature sequences and receiver filters to
 - maximize the user capacity
 - maximize the information theoretic sum capacity
- Multicell case is an open problem. Assume a single-cell system from now on.

Controllable Parameters

- Each user has a power (p_i), a signature sequence (s_i) and a receiver filter (c_i) — single-cell



Information Theoretic CDMA Capacity

- Let $\mathbf{S} = [\mathbf{s}_1, \mathbf{s}_2, \dots, \mathbf{s}_K]$.
- C_{sum} : maximum total number of bits K users can transmit on the uplink (Verdú, 1986)

$$C_{\text{sum}} = \frac{1}{2} \log \left[\det \left(\mathbf{I}_N + \sigma^{-2} \mathbf{S} \mathbf{P} \mathbf{S}^\top \right) \right]$$

- For equal powers $p_i = p$ for all i ,

$$C_{\text{sum}} = \frac{1}{2} \log \left[\det \left(\mathbf{I}_N + \frac{p}{\sigma^2} \mathbf{S} \mathbf{S}^\top \right) \right] = \frac{1}{2} \log \left[\det \left(\mathbf{I}_K + \frac{p}{\sigma^2} \mathbf{S}^\top \mathbf{S} \right) \right]$$

- To maximize the sum capacity (Rupf-Massey, 1994)
 - If $K \leq N$, $\mathbf{S}^\top \mathbf{S} = \mathbf{I}_K$ K orthonormal sequences
 - If $K > N$, $\mathbf{S} \mathbf{S}^\top = \frac{K}{N} \mathbf{I}_N$ K Welch Bound Equality sequences

User Capacity

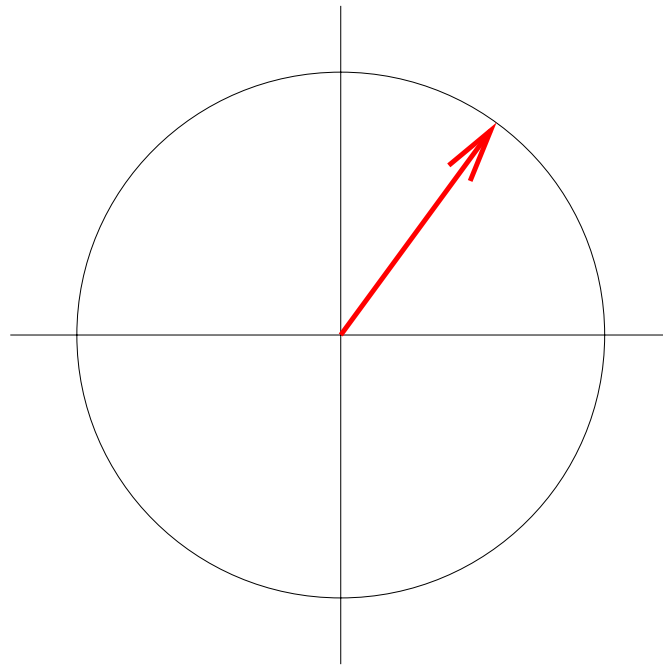
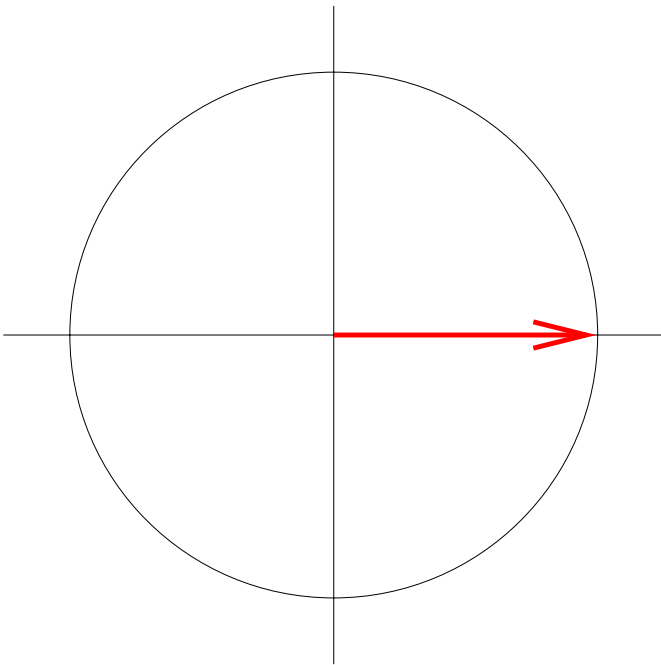
- Maximum number of admissible users given processing gain N and SIR target β
- K users are *admissible* if there are $p_i \geq 0$ and $\mathbf{s}_i, \mathbf{c}_i \in \mathbf{R}^N$ such that $\text{SIR}_i \geq \beta$, for all i
- User capacity of single-cell synchronous CDMA (Viswanath-Anantharam-Tse, 1999)

$$K < N(1 + 1/\beta)$$

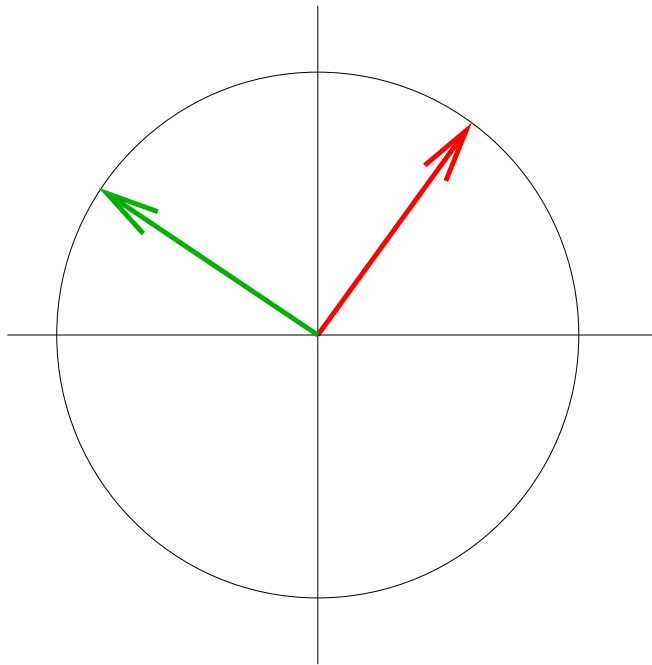
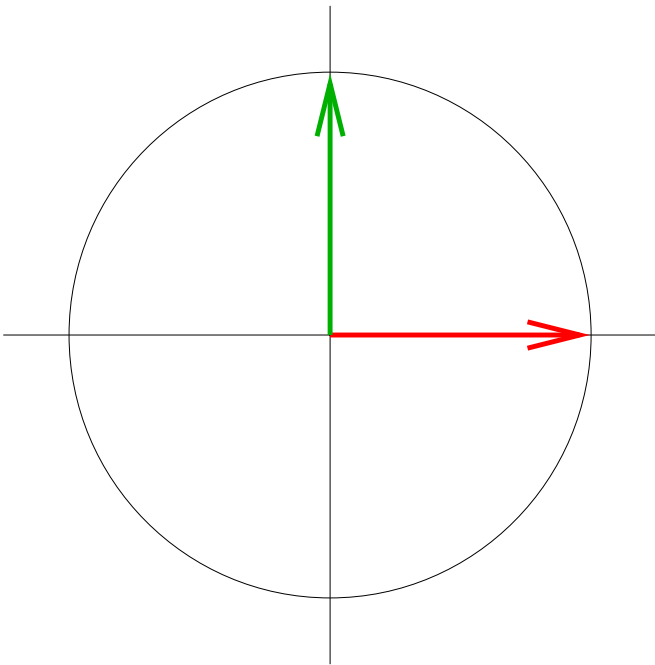
$$\beta < \frac{1}{\frac{K}{N} - 1}$$

- Optimum deterministic configuration
 - Equal received powers: $p_i = p$ for all i
 - WBE signature sequences: $\mathbf{S}\mathbf{S}^\top = (K/N)\mathbf{I}_N$
- For this configuration: MMSE receiver filters are scaled matched filters!

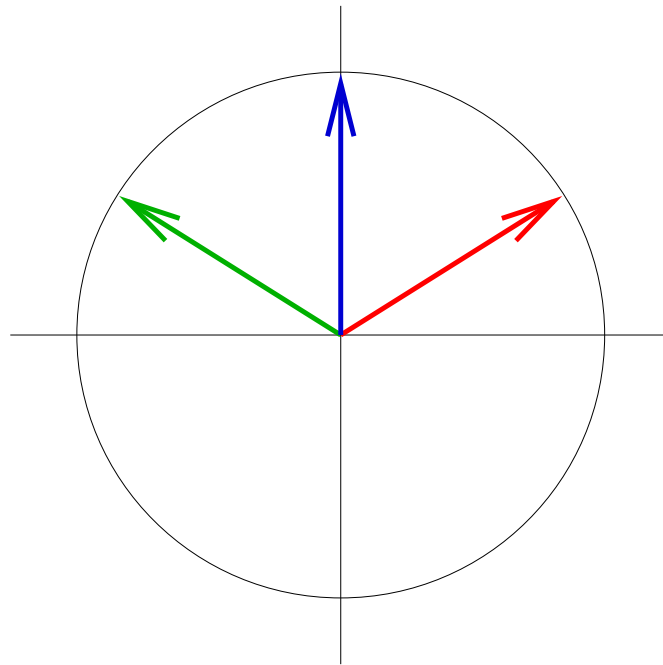
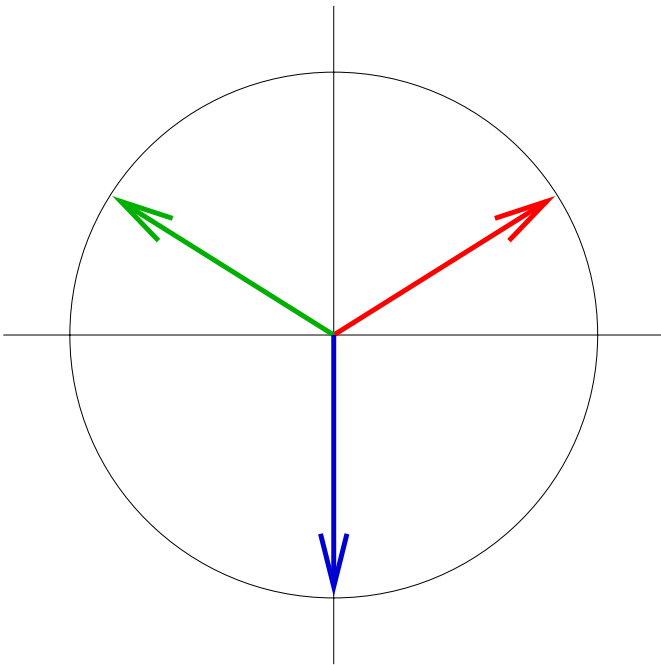
Simple example, $N = 2, K = 1$



Simple example, $N = 2, K = 2$



Simple example, $N = 2, K = 3$



Welch's Bound, Minimum TSC and Optimality

- For K unit-energy vectors $\{\mathbf{s}_1, \dots, \mathbf{s}_K\}$, Total Squared Correlation (TSC) (Welch, 1974):

$$\text{TSC} = \sum_{i=1}^K \sum_{j=1}^K (\mathbf{s}_i^\top \mathbf{s}_j)^2 \geq \frac{K^2}{N}$$

- If $K \leq N$, the bound is loose: $\text{TSC} \geq K$. The bound is achieved iff $\mathbf{S}^\top \mathbf{S} = \mathbf{I}_K$.
- If $K > N$, the bound is achieved iff $\mathbf{S}\mathbf{S}^\top = \frac{K}{N}\mathbf{I}_N$ (Massey-Mittelholzer, 1991).
- **Minimum TSC \iff Optimum signature sequences**
- Goal: A simple algorithm to construct them.

TSC Reduction

- Separate terms that depend on \mathbf{s}_i in TSC

$$\text{TSC} = (\mathbf{s}_i^\top \mathbf{s}_i)^2 + 2\mathbf{s}_i^\top \left(\sum_{j \neq i} \mathbf{s}_j \mathbf{s}_j^\top \right) \mathbf{s}_i + \sum_{k \neq i} \sum_{l \neq i} (\mathbf{s}_k^\top \mathbf{s}_l)^2$$

- Many ways to reduce TSC.
- **MMSE update:** (Ulukus-Yates, 1998) replace \mathbf{s}_i with

$$\mathbf{c}_i = \alpha_i \mathbf{A}_i^{-1} \mathbf{s}_i$$

where $\mathbf{A}_i = \sum_{j \neq i} \mathbf{s}_j \mathbf{s}_j^\top + a^2 \mathbf{I}_N$, and α_i is the normalizing factor.

- \mathbf{c}_i is a generalized normalized MMSE filter for user i .
- **Eigen update:** (Rose-Ulukus-Yates, 2000) replace \mathbf{s}_i with the minimum eigenvector of \mathbf{A}_i .

Algorithm Properties

- Replace \mathbf{s}_i with normalized MMSE filter (or eigenvector) \mathbf{c}_i

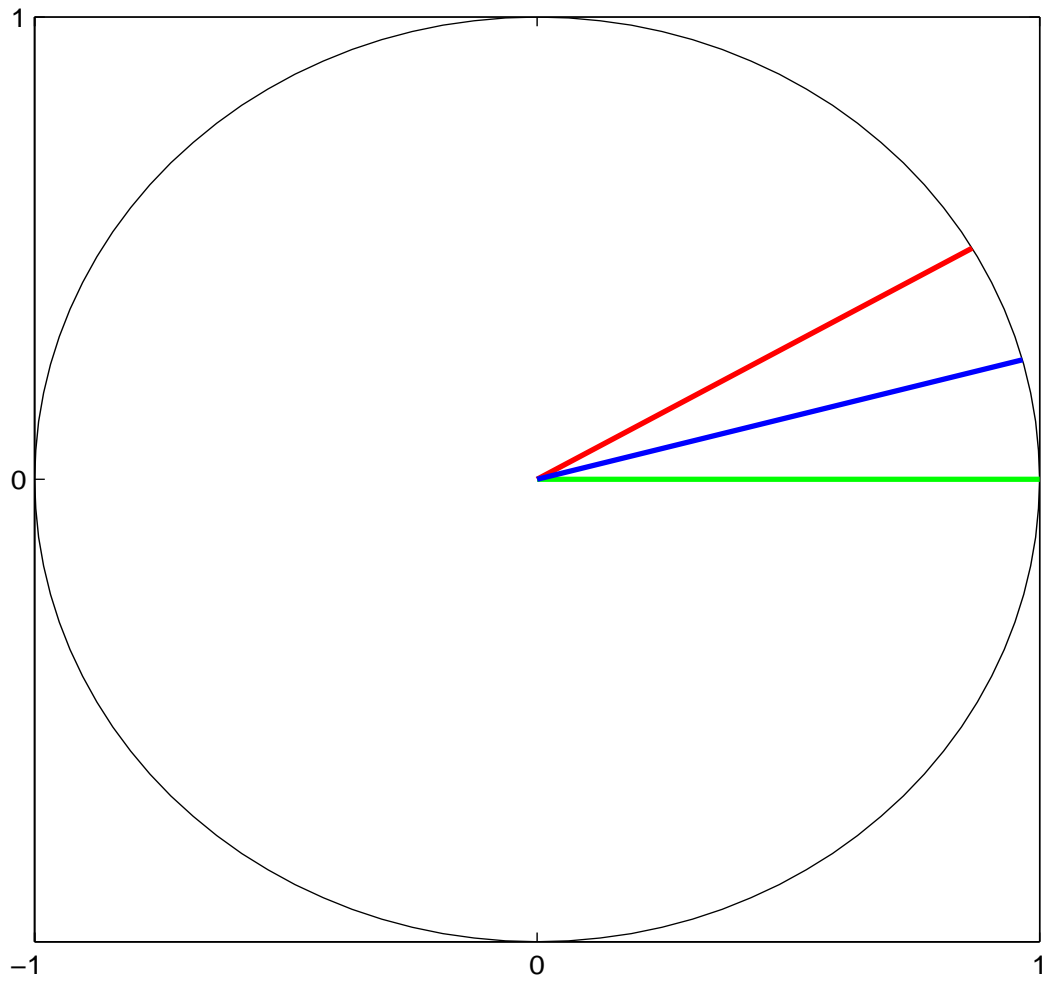
$$\mathbf{S} = [\mathbf{s}_1, \dots, \mathbf{s}_{i-1}, \mathbf{s}_i, \mathbf{s}_{i+1}, \dots, \mathbf{s}_K] \longrightarrow \mathbf{S}' = [\mathbf{s}_1, \dots, \mathbf{s}_{i-1}, \mathbf{c}_i, \mathbf{s}_{i+1}, \dots, \mathbf{s}_K]$$

- The proposed updates decrease the TSC, and increases the sum capacity:

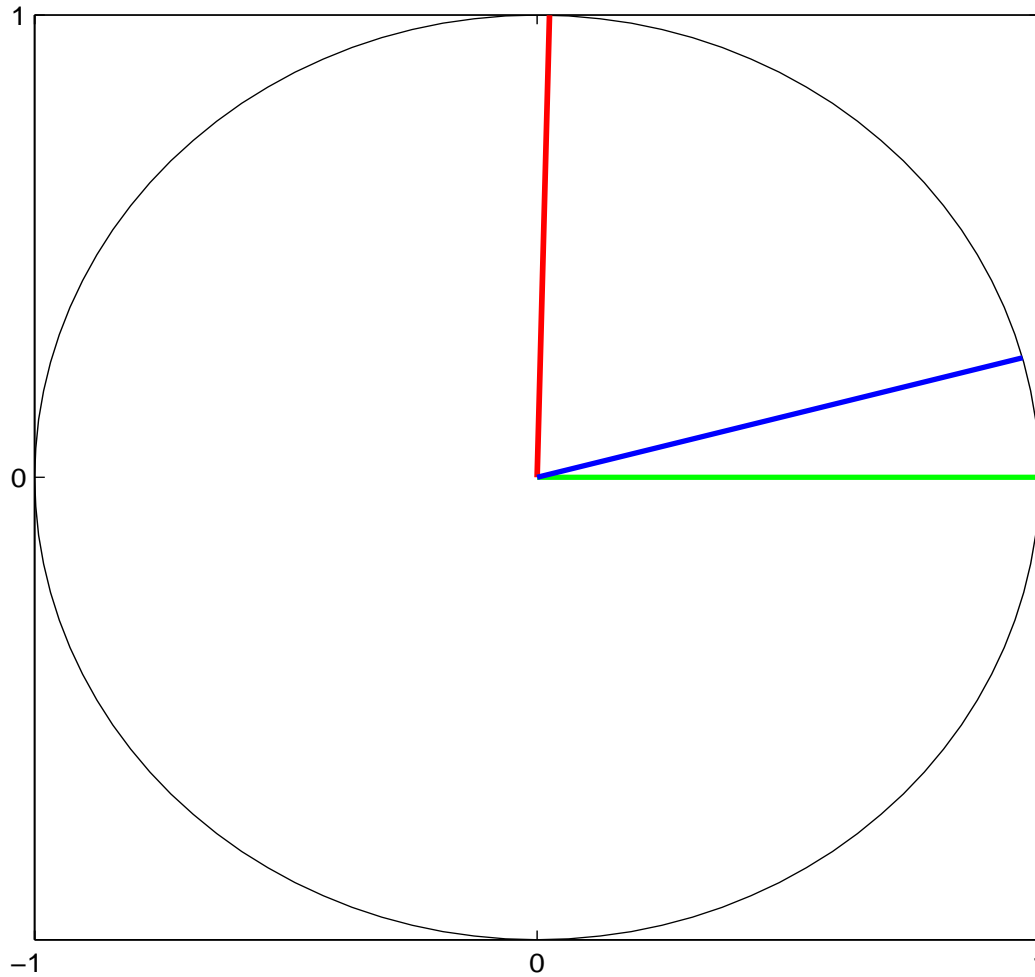
$$\begin{aligned} \text{TSC}(\mathbf{S}') &\leq \text{TSC}(\mathbf{S}) \\ C_{\text{sum}}(\mathbf{S}') &\geq C_{\text{sum}}(\mathbf{S}) \end{aligned}$$

- Let users update their signatures **sequentially** (one user at a time).
- The algorithm produces **progressively better** signature sequence sets and converges to an optimum set under mild conditions on the initial set.
- Need sequential updates condition for the convergence proof.

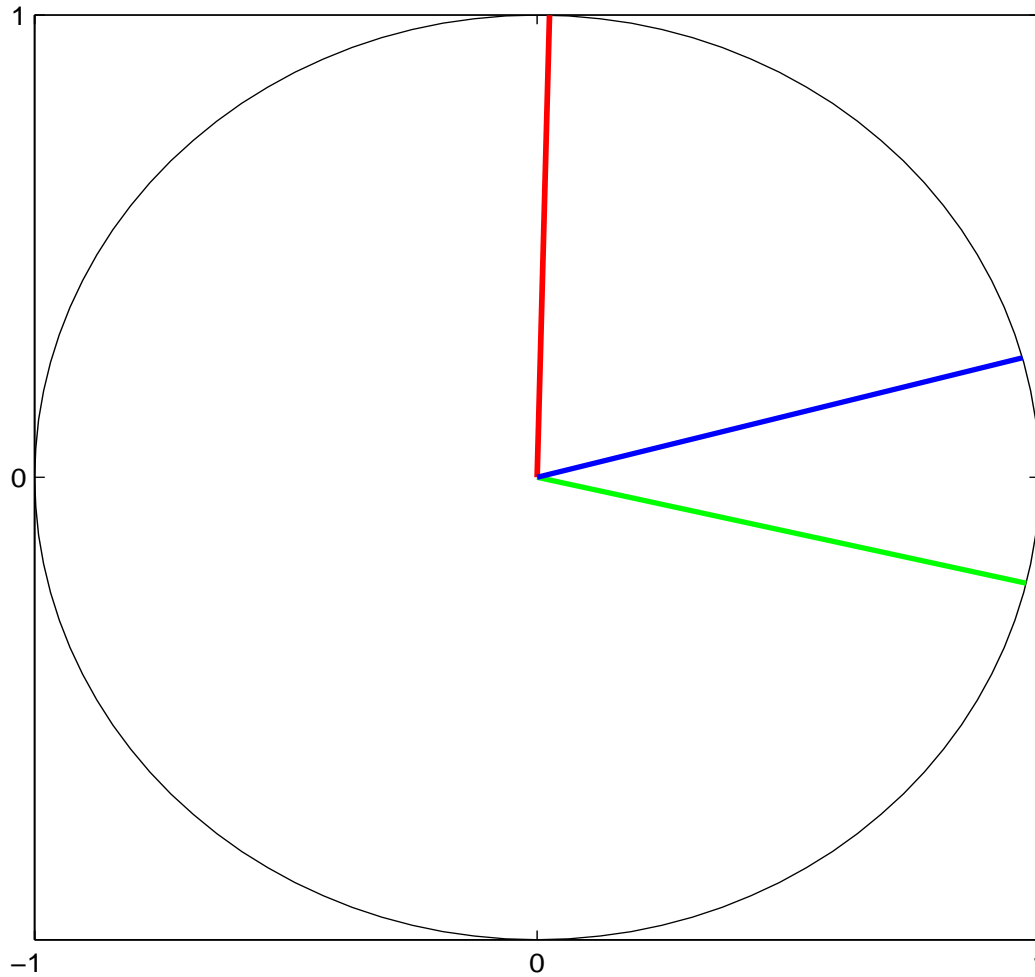
A Sample Run – start



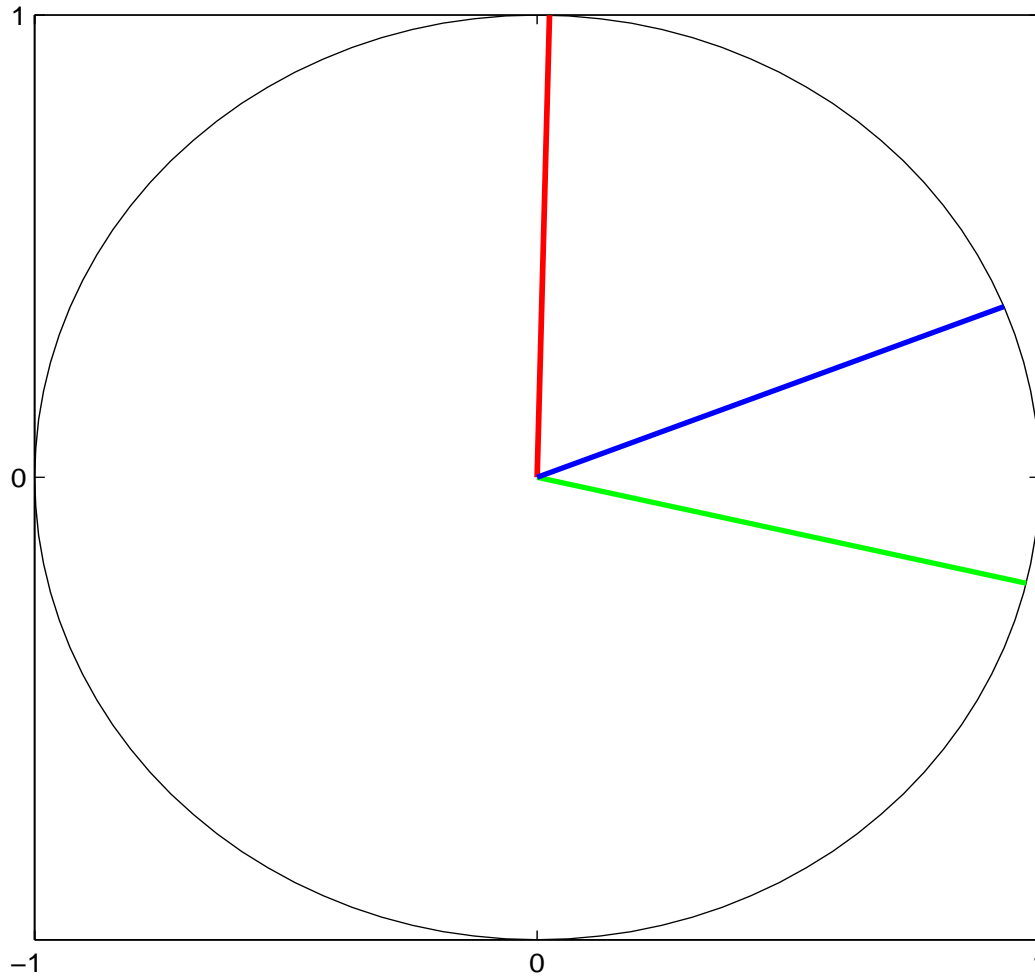
A Sample Run – 1st user updates



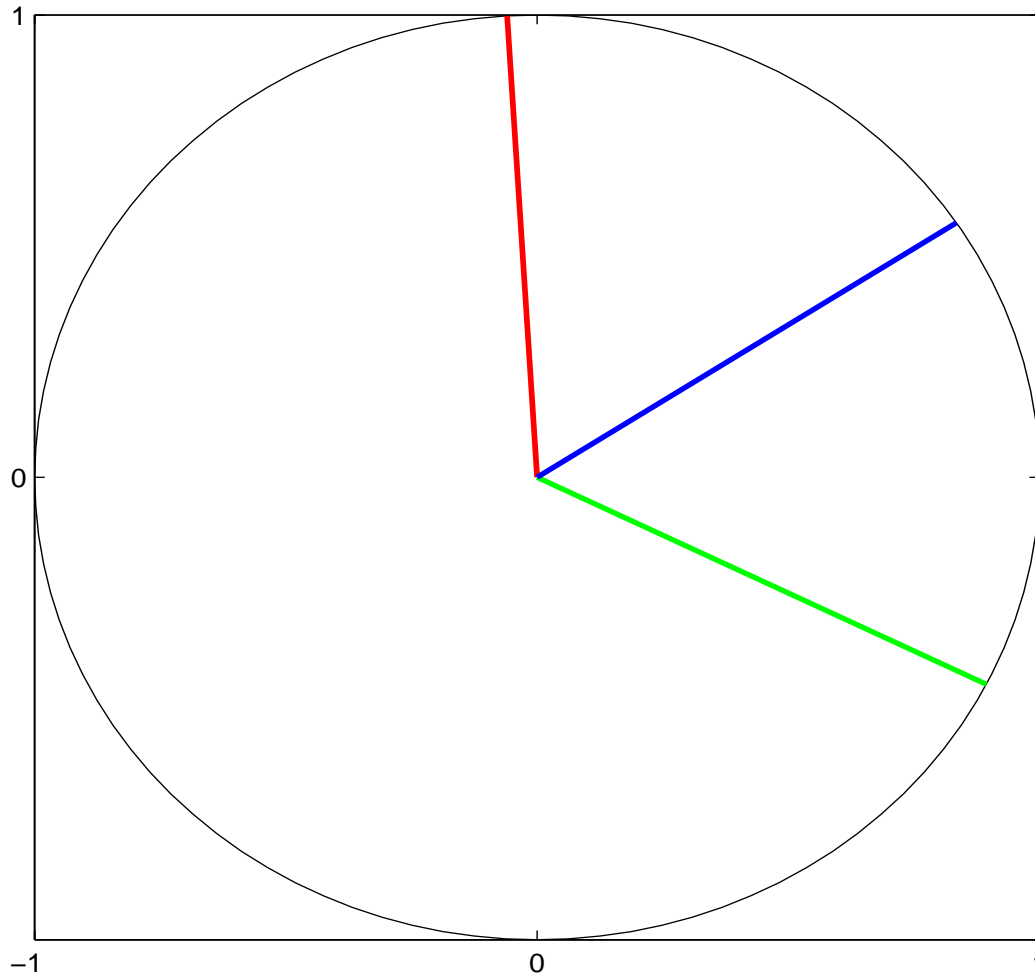
A Sample Run – 2nd user updates



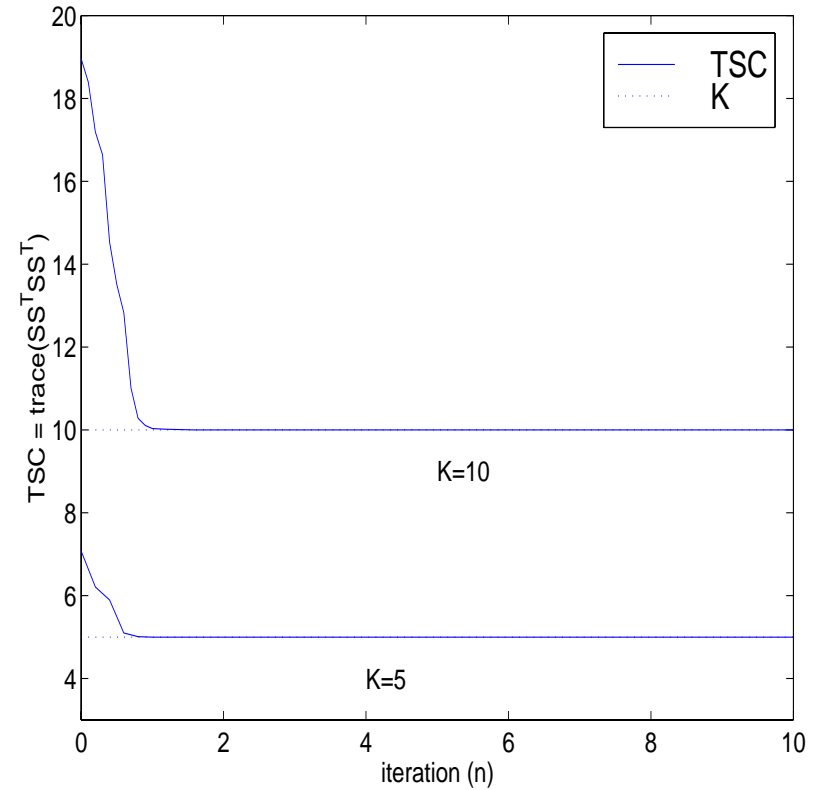
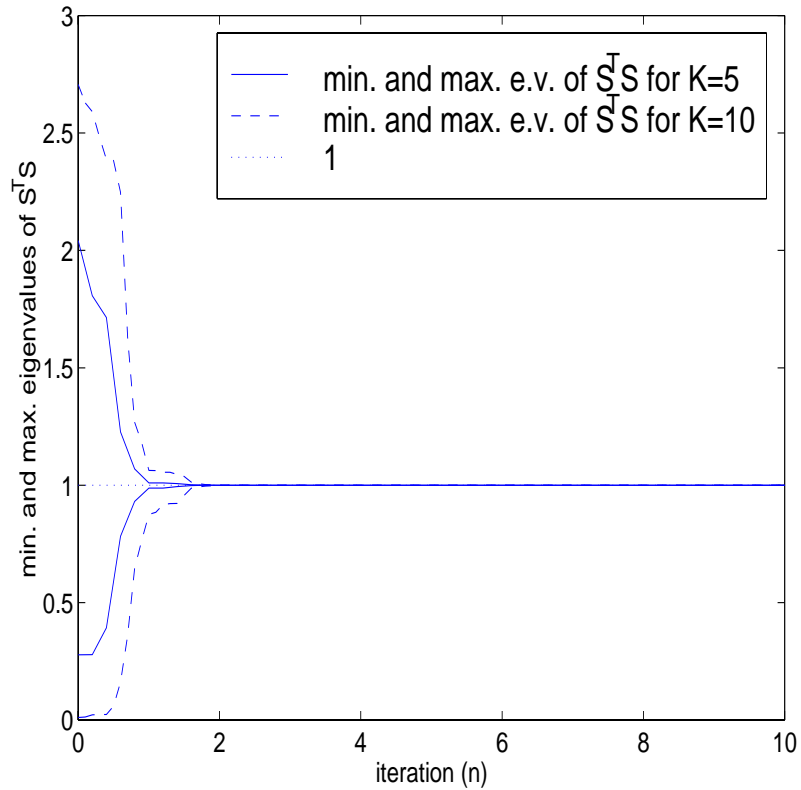
A Sample Run – 3rd user updates



A Sample Run – after convergence

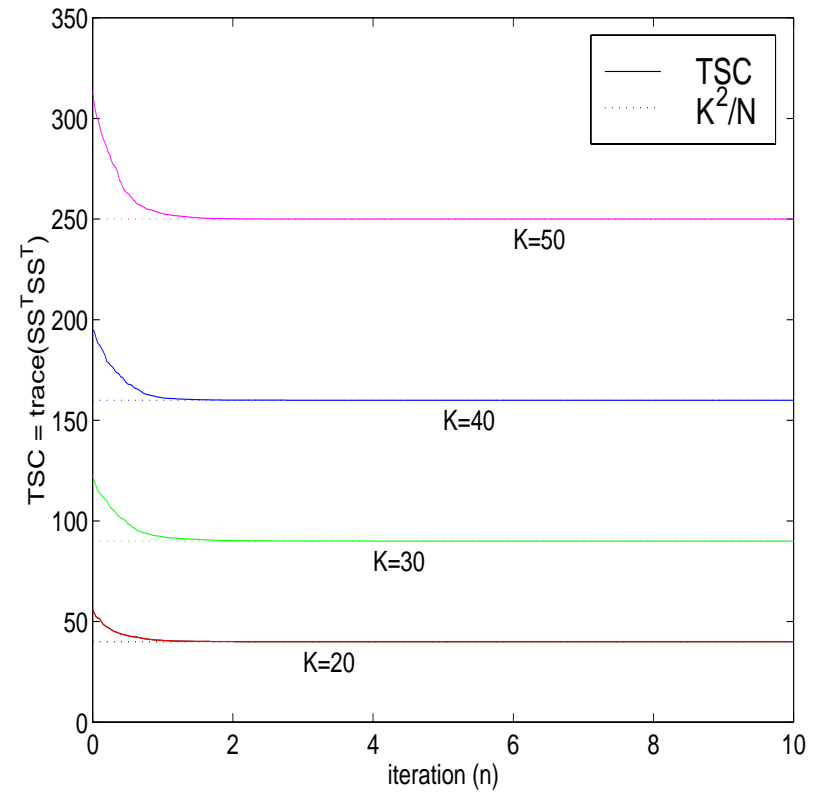
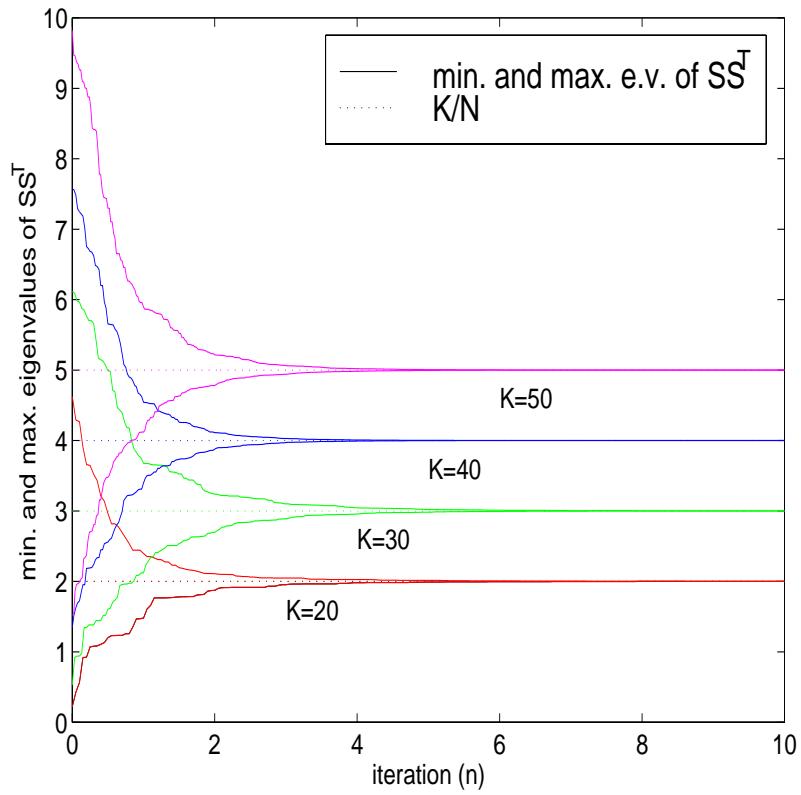


Min/Max Eigenvalues of $S^T S$ and TSC – MMSE Update



$K \leq N$ case: $K = 5, 10$ and $N = 10$

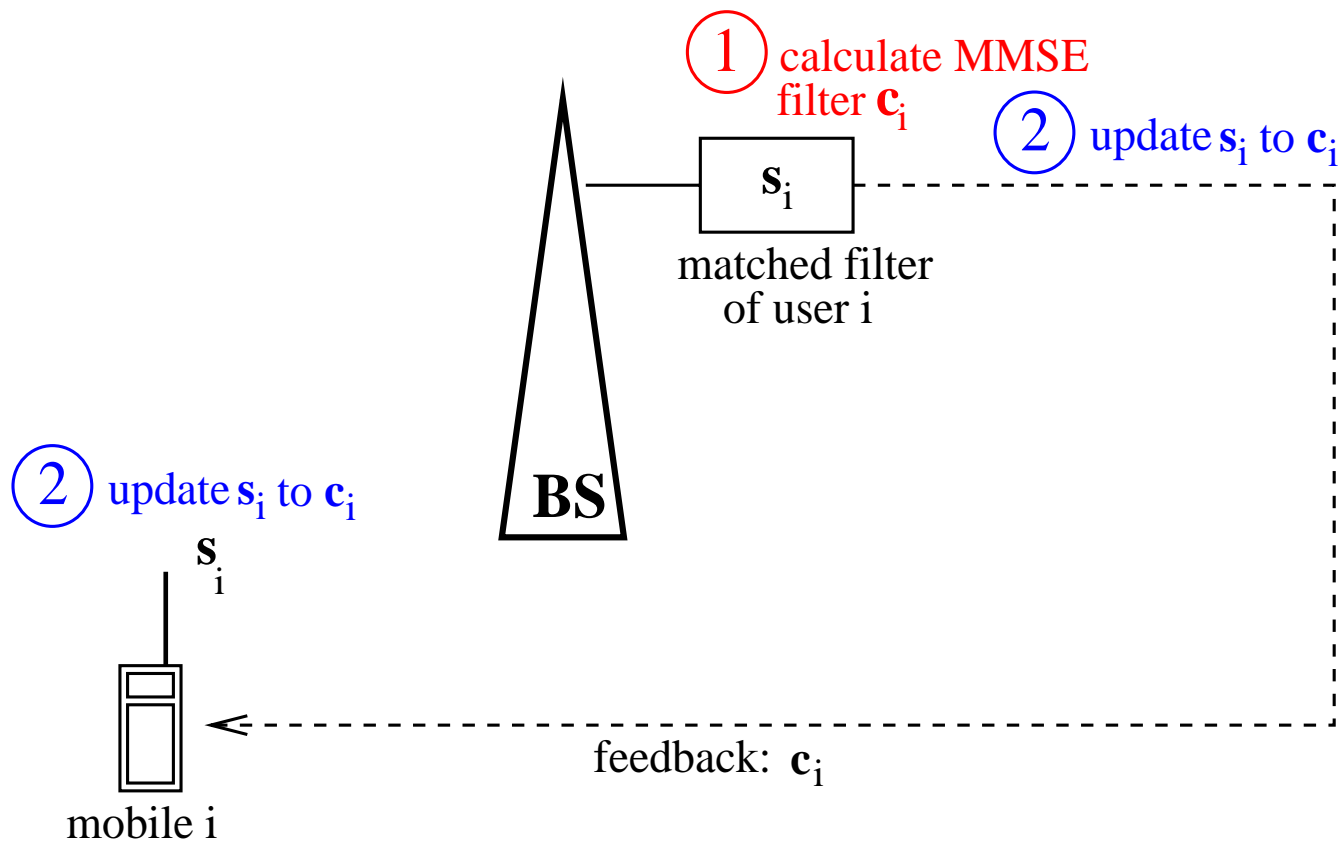
Min/Max Eigenvalues of SS^T and TSC – MMSE Update



$K > N$ case: $K = 20, 30, 40, 50$ and $N = 10$

Interference Avoidance – Implementation

- Calculate the receiver filter to be the MMSE filter
- Update the signature sequence and the receiver simultaneously

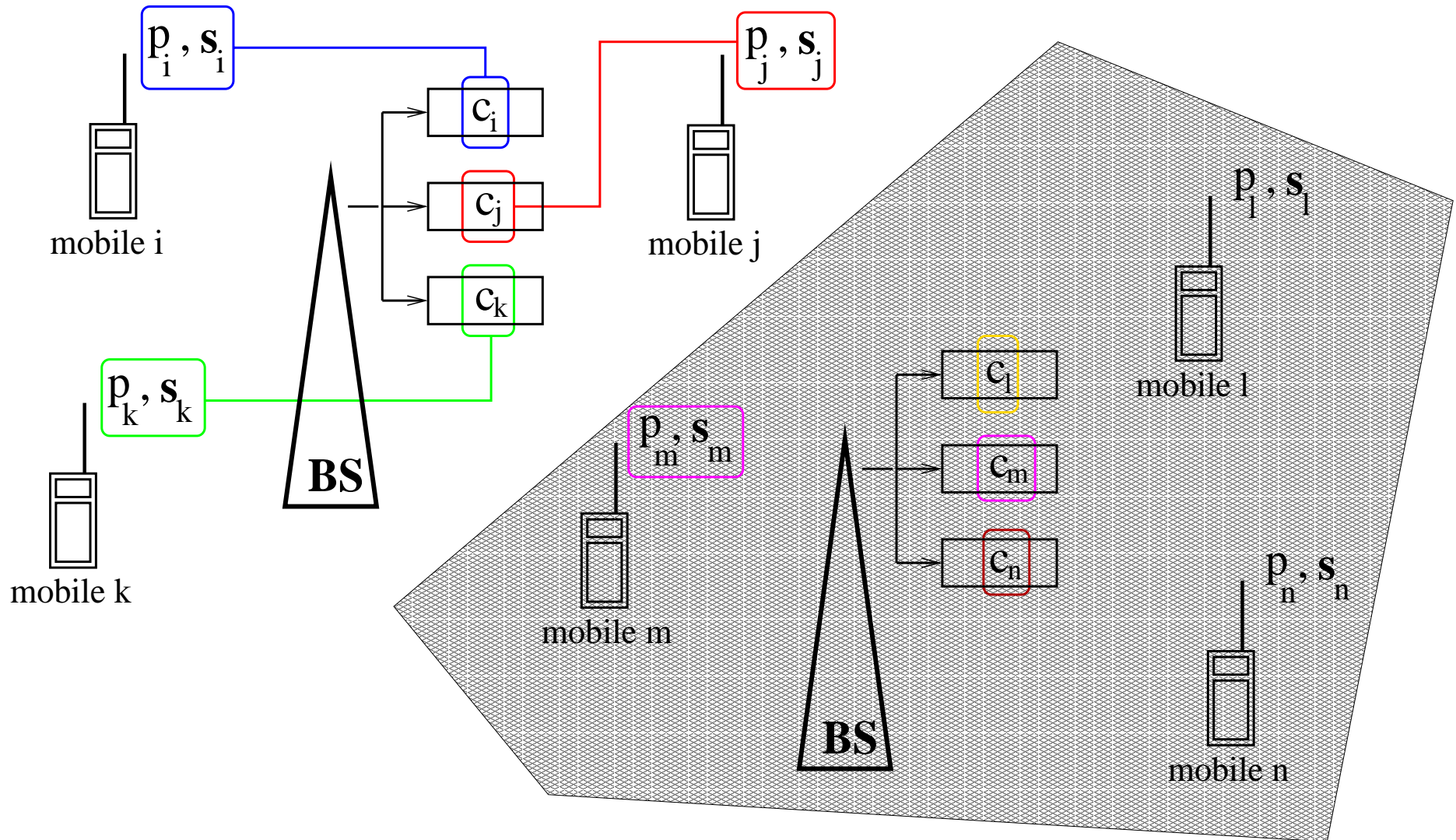


Further Improvements

- Possible shortcomings
 - Requires immediate feedback and **synchronization between transmitter and receiver**
 - Or, run signature adaptation algorithm **off-line**, when done, update transmitter once
 - For convergence proof, users need to take turns for signature updates, i.e., **sequential**
- Is an **on-line, parallel** algorithm **without synchronization** possible?
- Joint transmitter-receiver optimization (Uluks-Yener, 2002)

Controllable Parameters

- Each user has a power (p_i), a signature sequence (s_i) and a receiver filter (c_i) — single-cell



Related Performance Measures

- Need to introduce receiver filters into the optimization
- Remember C_{sum} and TSC are functions of signature sequences only
- Look at the MSE of user i , MSE_i

$$\text{MSE}_i = \mathbf{c}_i^\top \left(\mathbf{S}\mathbf{S}^\top + \sigma^2 \mathbf{I}_N \right) \mathbf{c}_i - 2\mathbf{c}_i^\top \mathbf{s}_i + 1$$

- Define the total MSE

$$\text{MSE} = \sum_{i=1}^K \text{MSE}_i = \text{tr} \left[\mathbf{C}^\top \left(\mathbf{S}\mathbf{S}^\top + \sigma^2 \mathbf{I}_N \right) \mathbf{C} - 2\mathbf{C}^\top \mathbf{S} + \mathbf{I}_K \right]$$

where $\mathbf{C} = [\mathbf{c}_1, \dots, \mathbf{c}_K]$.

- **MSE is a function of signature sequences and receiver filters.**

Equivalent Optimization Problems

- Using MMSE filters as receivers, we get the system-wide MMSE

$$\text{MMSE} = K - \text{tr} \left[\mathbf{S}\mathbf{S}^\top \left(\mathbf{S}\mathbf{S}^\top + \sigma^2 \mathbf{I}_N \right)^{-1} \right] = K - \sum_{i=1}^N \frac{\lambda_i}{\lambda_i + \sigma^2}$$

$$\text{TSC} = \text{tr} \left[\left(\mathbf{S}\mathbf{S}^\top \right)^2 \right] = \sum_{i=1}^N \lambda_i^2$$

$$C_{\text{sum}} = \frac{1}{2} \log \left[\det \left(\mathbf{I}_N + \sigma^{-2} \mathbf{S}\mathbf{S}^\top \right) \right] = \frac{1}{2} \sum_{i=1}^N \log \left(1 + \frac{\lambda_i}{\sigma^2} \right)$$

- C_{sum} is **Schur-concave**, and MMSE and TSC are **Schur-convex** functions of $\{\lambda_i\}_{i=1}^N$.
- **Signature set that maximizes C_{sum} minimizes MMSE and TSC, and vice versa.**

MSE Minimization

- Minimize MSE over receivers and unit energy signatures

$$\begin{aligned} \min_{\{\mathbf{c}_i, \mathbf{s}_i\}} \quad & \text{tr} \left[\mathbf{C}^\top \left(\mathbf{S}\mathbf{S}^\top + \sigma^2 \mathbf{I}_N \right) \mathbf{C} - 2\mathbf{C}^\top \mathbf{S} + \mathbf{I}_K \right] \\ \text{s.t.} \quad & \mathbf{s}_i^\top \mathbf{s}_i = 1 \quad i = 1, \dots, K \end{aligned}$$

- An iterative algorithm using **alternating minimization**
- Receiver update for user i :

$$\mathbf{c}_i = \left(\mathbf{S}\mathbf{S}^\top + \sigma^2 \mathbf{I}_N \right)^{-1} \mathbf{s}_i$$

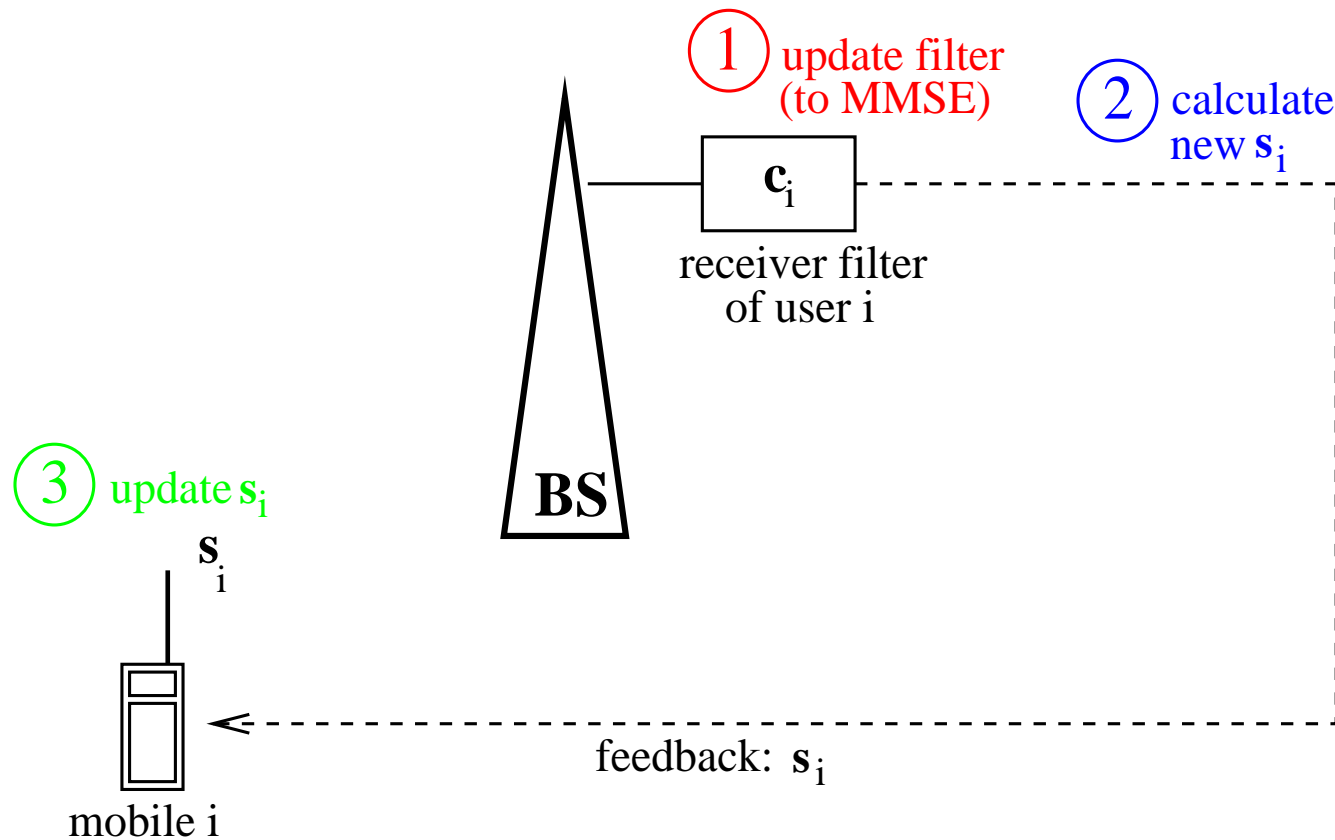
- Users can update their receivers independently, i.e., **in parallel**
- Transmitter update for user i :

$$\mathbf{s}_i = \left(\mathbf{C}\mathbf{C}^\top + \alpha_i \mathbf{I}_N \right)^{-1} \mathbf{c}_i$$

- Users can update transmitters independently, i.e., **in parallel**

Joint Transmitter/Receiver Optimization – Implementation

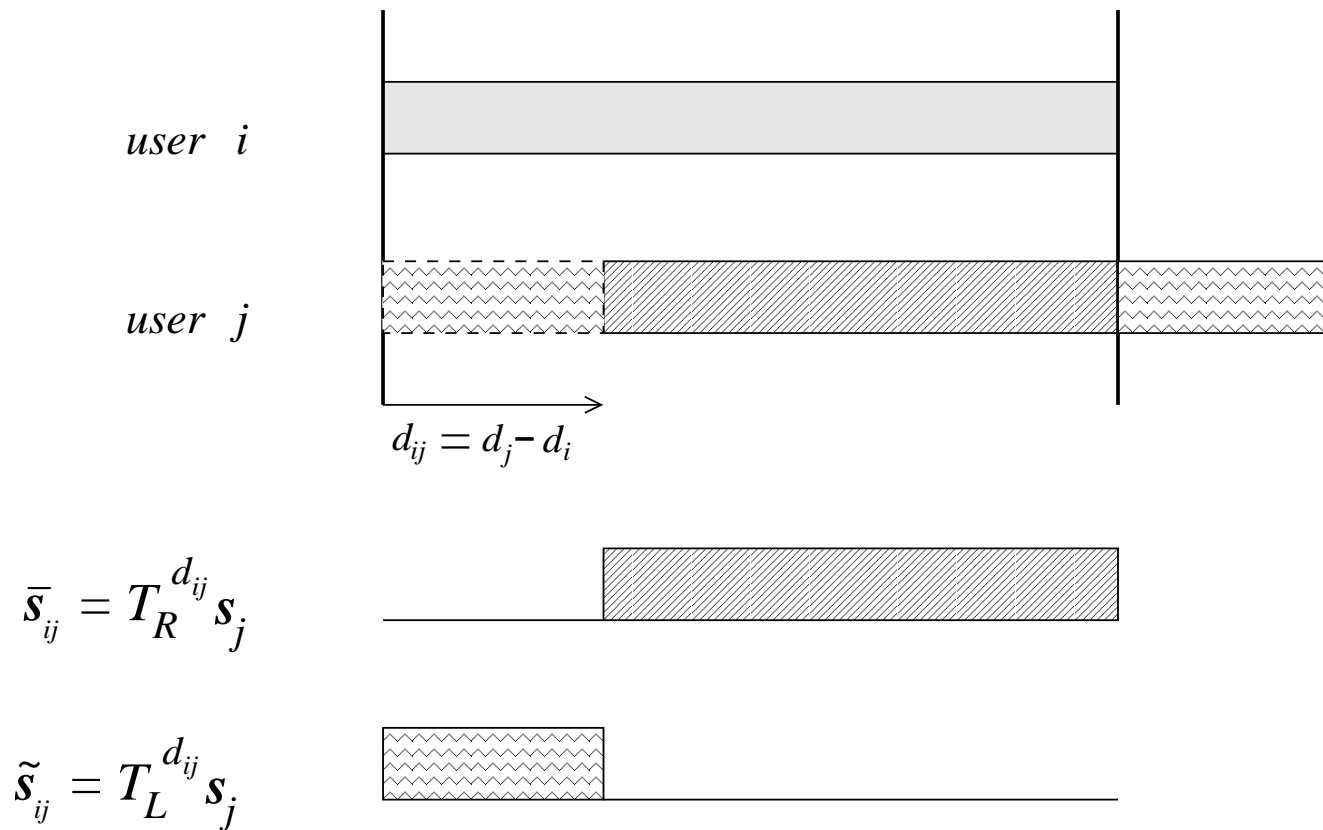
- Update the receiver filter to be the MMSE filter
- Calculate new signature sequence with the updated filter
- Update the signature sequence



User Capacity in Asynchronous System

- What is the user capacity of a single-cell asynchronous CDMA system?
- How do we choose signature sequences? Optimality criteria?
- Traditionally: choose signatures with *uniformly good* auto and cross correlations for all delay profiles [Pursley, Pursley-Sarwate, Massey-Uhran, Mow].
- In every wireless system, for some time, the delays of the users are fixed
- **Design signature sequences that are *best* for the current fixed delay profile**
- without having to know the current delay profile!

Asynchronous Interference Representation



- Received signal of the i th user in one symbol interval

$$\mathbf{r}_i = \sqrt{p_i} b_i \mathbf{s}_i + \sum_{j \neq i} \sqrt{p_j} (\bar{b}_j \bar{\mathbf{s}}_{ij} + \tilde{b}_j \tilde{\mathbf{s}}_{ij}) + \mathbf{n}_i$$

Asynchronous cross-correlations

- Define total squared asynchronous correlation (Uluks-Yates, 2000)

$$\text{TSAC} = \sum_{i=1}^K \sum_{j=1}^K \left\{ (\mathbf{s}_i^\top \bar{\mathbf{s}}_{ij})^2 + (\mathbf{s}_i^\top \tilde{\mathbf{s}}_{ij})^2 \right\}$$

- Optimum signature sequences minimize the TSAC.
- Optimum powers are all equal.
- The user capacity of the asynchronous system is the same as the synchronous system.
- **There is no loss in user-capacity due to asynchrony.**
- With optimum signature sequences, single-user, one-shot matched filters are optimum.

TSAC Reduction

- Optimum signature sequences minimize the TSAC
- Similar to the synchronous case, separate the terms that depend on \mathbf{s}_i

$$\text{TSAC} = (\mathbf{s}_i^\top \mathbf{s}_i)^2 + 2\mathbf{s}_i^\top \left(\sum_{j \neq i} \bar{\mathbf{s}}_{ij} \bar{\mathbf{s}}_{ij}^\top + \tilde{\mathbf{s}}_{ij} \tilde{\mathbf{s}}_{ij}^\top \right) \mathbf{s}_i + \sum_{k \neq i} \sum_{l \neq i} R_{kl}$$

- Many TSAC reduction algorithms are possible
- **Asynchronous MMSE update:**

$$\mathbf{c}_i = \alpha_i \mathbf{A}_i^{-1} \mathbf{s}_i$$

where $\mathbf{A}_i = \sum_{j \neq i} \bar{\mathbf{s}}_{ij} \bar{\mathbf{s}}_{ij}^\top + \tilde{\mathbf{s}}_{ij} \tilde{\mathbf{s}}_{ij}^\top + a^2 \mathbf{I}_N$.

- **Asynchronous eigen update:**

$\mathbf{c}_i =$ normalized eigenvector of \mathbf{A}_i with the smallest eigenvalue

Some Open Problems

- Multipath channel without ISI (some initial results available)
- Multipath channel with ISI
- Fading channels
- Channels with multiple transmit and/or multiple receive antennas
 - interactions between spatial and temporal signatures
- Multicell channel
- Signal design for security:
 - definition of secure signature sequences
- Signal design for cross-layer interactions:
 - e.g., jointly optimum signatures and route/flow assignments