Introduction to Cryptology

Lecture 7

Announcements

- HW3 due Tuesday, 2/21
- Quiz Solutions up on Canvas

Agenda

- Last time:
  - The Computational Approach (K/L 3.1)
  - Defining computationally secure SKE (K/L 3.2)
- This time:
  - Defining PRG (K/L 3.3)
  - Exercise on PRG
  - Constructing computationally secure SKE (K/L 3.3)
  - Security proof for construction (K/L 3.3)
  - Discussion on Stream Ciphers

Pseudorandom Generator

- Functionality
  - Deterministic algorithm G
  - Takes as input a short random seed s
  - Ouputs a long string $G(s)$
- Security
  - No efficient algorithm can “distinguish” $G(s)$ from a truly random string r.
  - i.e. passes all “statistical tests.”
- Intuition:
  - Stretches a small amount of true randomness to a larger amount of pseudorandomness.
- Why is this useful?
  - We will see that pseudorandom generators will allow us to beat the Shannon bound of $|\boldsymbol{K}| \geq |\boldsymbol{M}|$.
  - I.e. we will build a computationally secure encryption scheme with $|\boldsymbol{K}| < |\boldsymbol{M}|$

Pseudorandom Generators

Definition: Let $\ell(\cdot)$ be a polynomial and let G be a deterministic poly-time algorithm such that for any input $s \in \{0,1\}^n$, algorithm G outputs a string of length $\ell(n)$. We say that G is a pseudorandom generator if the following two conditions hold:

1. (Expansion:) For every n it holds that $\ell(n) > n$.
2. (Pseudorandomness:) For all ppt distinguishers D, there exists a negligible function $negl$ such that:

$$\left|\Pr[D(r) = 1] - \Pr\left[D\left(G(s)\right) = 1\right]\right| \leq negl(n),$$

where r is chosen uniformly at random from $\{0,1\}^{\ell(n)}$, the seed s is chosen uniformly at random from $\{0,1\}^n$, and the probabilities are taken over the random coins used by D and the choice of r and s.

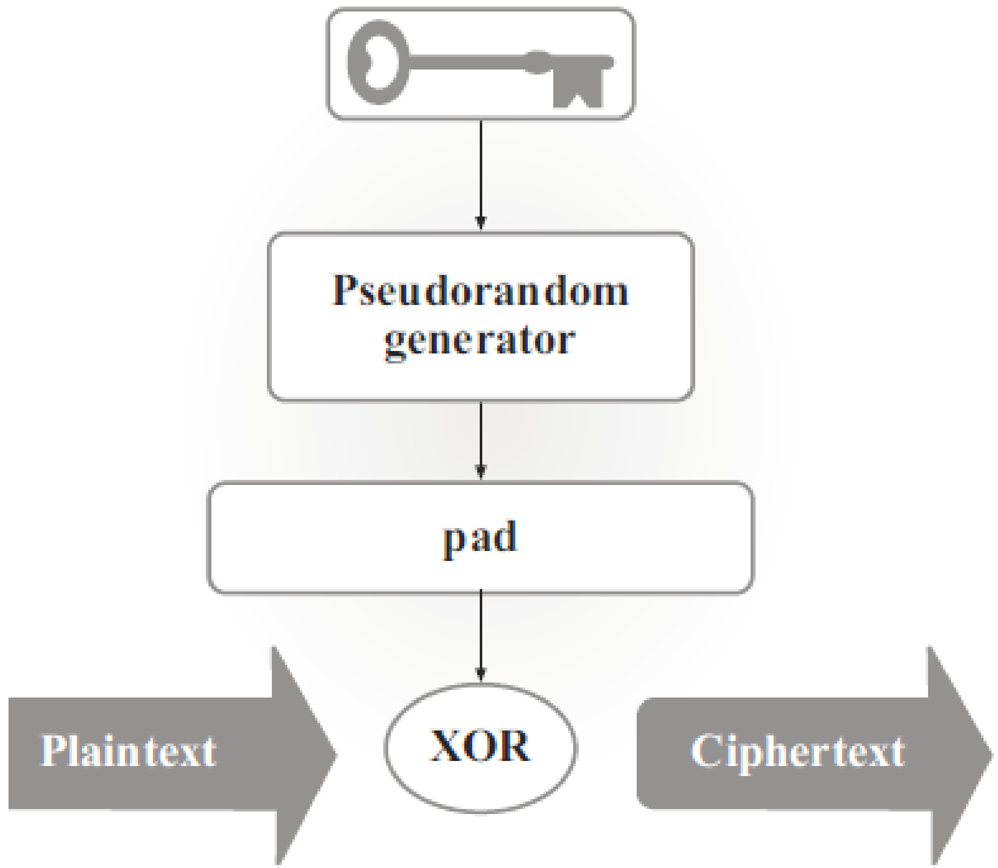
The function $\ell(\cdot)$ is called the expansion factor of G.

Stream Cipher

- Practical instantiation of a pseudorandom generator (will talk more about them and how they are constructed later in the course).
- Pseudorandom bits of a stream cipher are produced gradually and on demand.
- Application can request exact number of bits needed.
- This improves efficiency.

Constructing Secure Encryption Schemes

A Secure Fixed-Length Encryption Scheme



The Encryption Scheme

Let G be a pseudorandom generator with expansion factor ℓ. Define a private-key encryption scheme for messages of length ℓ as follows:

- Gen: on input 1^n, choose $k \leftarrow \{0,1\}^n$ uniformly at random and output it as the key.
- Enc: on input a key $k \in \{0,1\}^n$ and a message $m \in \{0,1\}^{\ell(n)}$, output the ciphertext
$$c := G(k) \oplus m.$$
- Dec: on input a key $k \in \{0,1\}^n$ and a ciphertext $c \in \{0,1\}^{\ell(n)}$, output the plaintext message
$$m := G(k) \oplus c.$$

Recall: Indistinguishability in the presence of an eavesdropper

Consider a private-key encryption scheme $\Pi = (Gen, Enc, Dec)$, any adversary A, and any value n for the security parameter.

The eavesdropping indistinguishability experiment $PrivK^{eav}{}_{A,\Pi}(n)$:

1. The adversary A is given input 1^n, and outputs a pair of messages m_0, m_1 of the same length.
2. A key k is generated by running $Gen(1^n)$, and a random bit $b \leftarrow \{0,1\}$ is chosen. A challenge ciphertext $c \leftarrow Enc_k(m_b)$ is computed and given to A.
3. Adversary A outputs a bit b'.
4. The output of the experiment is defined to be 1 if $b' = b$, and 0 otherwise. If $PrivK^{eav}{}_{A,\Pi}(n) = 1$, we say that A succeeded.

Recall: Indistinguishability in the presence of an eavesdropper

Definition: A private key encryption scheme $\Pi = (Gen, Enc, Dec)$ has indistinguishable encryptions in the presence of an eavesdropper if for all probabilistic polynomial-time adversaries A there exists a negligible function $negl$ such that

$$\Pr\left[PrivK^{eav}{}_{A,\Pi}(n) = 1\right] \leq \frac{1}{2} + negl(n),$$

Where the prob. Is taken over the random coins used by A, as well as the random coins used in the experiment.

Security Analysis

Theorem: If G is a pseudorandom generator, then the Construction above is a fixed-length private-key encryption scheme that has indistinguishable encryptions in the presence of an eavesdropper.

Security Analysis

- Proof by reduction method.

Security Analysis

Proof: Let A be a ppt adversary trying to break the security of the construction. We construct a distinguisher D that uses A as a subroutine to break the security of the PRG.

Distinguisher D:

D is given as input a string $w \in \{0,1\}^{\ell(n)}$.

1. Run $A(1^n)$ to obtain messages $m_0, m_1 \in \{0,1\}^{\ell(n)}$.
2. Choose a uniform bit $b \in \{0,1\}$. Set $c \coloneqq w \oplus m_b$.
3. Give c to A and obtain output b'. Output 1 if $b' = b$, and output 0 otherwise.

Security Analysis

Consider the probability D outputs 1 in the case that w is random string r vs. w is a pseudorandom string $G(s)$.

- When w is random, D outputs 1 with probability exactly ½. Why?
- When w is pseudorandom, D outputs 1 with probability $\Pr\left[PrivK^{eav}{}_{A,\Pi}(n) = 1\right] = \frac{1}{2} + \rho(n)$, where ρ is non-negligible.

Security Analysis

D's distinguishing probability is:

$$\left|\frac{1}{2} - \left(\frac{1}{2} + \rho(n)\right)\right| = \rho(n).$$

This is a contradiction to the security of the PRG, since ρ is non-negligible.