

Cryptography

Lecture 18

Announcements

- HW6 due 4/13
- HW7 is up on the course webpage, due on 4/22

Agenda

- New Unit: Number Theory!

Modular Arithmetic

Definition of modulo:

We say that two integers a, b are congruent modulo p denoted by

$$a \equiv b \pmod{p}$$

If

$$p \mid (a - b)$$

(i.e. p divides $(a - b)$).

Modular Arithmetic

Examples: All of the following are true

$$2 \equiv 15 \pmod{13}$$

$$28 \equiv 15 \pmod{13}$$

$$41 \equiv 15 \pmod{13}$$

$$-11 \equiv 15 \pmod{13}$$

Modular Arithmetic

Operation: addition mod p

Regular addition, take modulo p .

Example: $8 + 10 \text{ mod } 13 \equiv 18 \text{ mod } 13 \equiv 5 \text{ mod } 13$.

Properties of Addition mod p

Consider the set Z_p of integers $\{0, 1, \dots, p - 1\}$ and the operation addition mod p .

- Closure: Adding two numbers in Z_p and taking mod p yields a number in Z_p .
- Identity: For every $a \in Z_p$, $[0 + a] \text{ mod } p \equiv a \text{ mod } p$.
- Inverse: For every $a \in Z_p$, there exists a $b \in Z_p$ such that $a + b \equiv 0 \text{ mod } p$.
 - b is simply the negation of a ($b = -a$).
 - Note that using the property of inverse, we can do subtraction. We define $c - d \text{ mod } p$ to be equivalent to $c + (-d) \text{ mod } p$.
- Associativity: For every $a, b, c \in Z_p$: $(a + b) + c = a + (b + c) \text{ mod } p$.

Z_p is a group with respect to addition!

Definition of a Group

A group is a set G along with a binary operation \circ for which the following conditions hold:

- Closure: For all $g, h \in G$, $g \circ h \in G$.
- Identity: There exists an identity $e \in G$ such that for all $g \in G$, $e \circ g = g = g \circ e$.
- Inverse: For all $g \in G$ there exists an element $h \in G$ such that $g \circ h = e = h \circ g$. Such an h is called an inverse of g .
- Associativity: For all $g_1, g_2, g_3 \in G$, $(g_1 \circ g_2) \circ g_3 = g_1 \circ (g_2 \circ g_3)$.

When G has a finite number of elements, we say G is finite and let $|G|$ denote the order of the group.

Abelian Group

A group G with operation \circ is abelian if the following holds:

- Commutativity: For all $g, h \in G$, $g \circ h = h \circ g$.

We will always deal with finite, abelian groups.

Other groups over the integers

- We will be interested mainly in multiplicative groups over the integers, since there are computational problems believed to be hard over such groups.
 - Such hard problems are the basis of number-theoretic cryptography.
- Group operation is multiplication mod p , instead of addition mod p .

Multiplication mod p

Example:

$$3 \cdot 8 \text{ mod } 13 \equiv 24 \text{ mod } 13 \equiv 11 \text{ mod } 13.$$

Multiplicative Groups

Is Z_p a group with respect to multiplication mod p ?

- Closure—YES
- Identity—YES (1 instead of 0)
- Associativity—YES
- Inverse—NO
 - 0 has no inverse since there is no integer a such that $0 \cdot a \equiv 1 \pmod{p}$.

Multiplicative Group

For p prime, define $Z_p^* = \{1, \dots, p - 1\}$ with operation multiplication mod p .

We will see that Z_p^* is indeed a multiplicative group!

To prove that Z_p^* is a multiplicative group, it is sufficient to prove that every element has a multiplicative inverse (since we have already argued that all other properties of a group are satisfied).

This is highly non-trivial, we will see how to prove it using the Euclidean Algorithm.

Inefficient method of finding inverses mod p

Example: Multiplicative inverse of 9 mod 11.

$$9 \cdot 1 \equiv 9 \pmod{11}$$

$$9 \cdot 2 \equiv 18 \equiv 7 \pmod{11}$$

$$9 \cdot 3 \equiv 27 \equiv 5 \pmod{11}$$

$$9 \cdot 4 \equiv 36 \equiv 3 \pmod{11}$$

$$9 \cdot 5 \equiv 45 \equiv 1 \pmod{11}$$

What is the time complexity?

Brute force search. In the worst case must try all 10 numbers in Z_{11}^* to find the inverse.

This is **exponential** time! Why? Inputs to the algorithm are (9,11). The length of the input is the length of the binary representation of (9,11). This means that input size is approx. $\log_2 11$ while the runtime is approx. $2^{\log_2 11} = 11$. The runtime is exponential in the input length.

Fortunately, there is an efficient algorithm for computing inverses.

Euclidean Algorithm

Theorem: Let a, p be positive integers. Then there exist integers X, Y such that $Xa + Yp = \gcd(a, p)$.

Given a, p , the Euclidean algorithm can be used to compute $\gcd(a, p)$ in polynomial time. The extended Euclidean algorithm can be used to compute X, Y in polynomial time.

Proving Z_p^* is a multiplicative group

In the following we prove that every element in Z_p^* has a multiplicative inverse when p is prime. This is sufficient to prove that Z_p^* is a multiplicative group.

Proof. Let $a \in Z_p^*$. Then $\gcd(a, p) = 1$, since p is prime.

By the Euclidean Algorithm, we can find integers X, Y such that $aX + pY = \gcd(a, p) = 1$.

Rearranging terms, we get that $pY = (aX - 1)$ and so $p \mid (aX - 1)$.

By definition of modulo, this implies that $aX \equiv 1 \pmod{p}$.

By definition of inverse, this implies that X is the multiplicative inverse of a .

Note: By above, the **extended Euclidean algorithm** gives us a way to **compute the multiplicative inverse in polynomial time**.

Extended Euclidean Algorithm

Example

Find: X, Y such that $9X + 23Y = \gcd(9, 23) = 1$.

$$23 = 2 \cdot 9 + 5$$

$$9 = 1 \cdot 5 + 4$$

$$5 = 1 \cdot 4 + 1$$

$$4 = 4 \cdot 1 + 0$$

$$1 = 5 - 1 \cdot 4$$

$$1 = 5 - 1 \cdot (9 - 1 \cdot 5)$$

$$1 = (23 - 2 \cdot 9) - (9 - (23 - 2 \cdot 9))$$

$$1 = 2 \cdot 23 - 5 \cdot 9$$

$-5 = 18 \pmod{23}$ is the multiplicative inverse of $9 \pmod{23}$.

Time Complexity of Euclidean Algorithm

When finding $\text{gcd}(a, b)$, the “ b ” value gets halved every two rounds.

Why?

Time complexity: $2\log(b)$.

This is polynomial in the length of the input.

Why?

Getting Back to Z_p^*

Group $Z_p^* = \{1, \dots, p - 1\}$ operation:
multiplication modulo p .

Order of a finite group is the number of
elements in the group.

Order of Z_p^* is $p - 1$.