ENEE/CMSC/MATH 456 Feistel Class Exercise

1. Consider a *two*-round Feistel Network with input length ℓ , key length n and round functions FF_{k}). where $F\{:0\}1^{n} \{ 0, \} \ell^{2} \to \{ 0, 1 \}^{n}$ is a pseudorandom function. Prove that the output of the Feistel Network is *not* a pseudorandom permutation (PRP).

See attached sheet for the structure of a Feistel Network.

Solution:

- 1. Query L_0||R_0, get back F_k(R_0) + L_0||* (we don't care what's on the right)
- 2. Query L_0 + \Delta || R_0, get back $F_k(R_0) + L_0 + Delta ||^*$
- 3. XOR the two left hand sides. If you get back \Delta, then output 0. Otherwise output 1.

The probability of outputting 1 when the oracle is a 2-round Feistel is 1. The probability of outputting 1 when the oracle is a random permutation is approximately 1/2^{\ell/2}. The difference is clearly non-negligible. Thus, this is a valid attack on the security of the 2-round Feistel.

2. ****Challenge**** Consider a *three*-round Feistel Network with input length ℓ , key length n and round functions $FFFF_k \cdot ,$ where $F \stackrel{\{: \ 0, 1 \ n \ \times \ 0, 1 \ \ell/2 \ \to \ 0, 1 \ \ell/2} \to 0, 1 \ \ell/2$ is a pseudorandom function. Prove that the output of the Feistel Network is *not* a strong pseudorandom permutation (sPRP).

See attached sheet for the structure of a Feistel Network.

Hint: The sequence of queries needed is:

- 1. Forward direction on (L_0||R_0), getting back (L_3||R_3)
- 2. Backward direction on $(L_3||R_3 + Delta)$, getting back $(L'_0||R'_0)$
- 3. Forward direction on (L_0 + \Delta||R_0), getting back (L''_3||R''_3)

There will be a relationship between R_0, L_3, R'_0 and L''_3

Solution: When we query $(L_0 || R_0)$, we get back $(L_3 || R_3)$.

When we query the backward direction on $(L_3||R_3+Delta)$, we get $L'_2 = R'_1 =$

 $L_2+Delta$, and $L'_1 = R'_0 = L_3 + f(L_2 + Delta)$, where $L_2 = L_0 + f(R_0)$

When we query the forward direction on (L_0+\Delta || R_0), we get back L''_3 = R_0 + f(L_0 + \Delta + f(R_0)) = R_0 + f(L_2 + \Delta).

So we have the relationship: $R'_0 + L_3 = R_0 + L''_3$.

If the above relationship occurs, the distinguisher outputs 1. The probability of outputting 1 when the oracle is 3-round Feistel is 1. The probability of outputting 1 when the oracle is a random permutation is approx. 1/2^{\ell/2}.



FIGURE 6.4: A 3-round Feistel network.